

Global properties of degree structures

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Overview

We will study structures from computability theory under the following aspects.

- Automorphism bases
- Biinterpretability with \mathbb{N} in params, prime model question
- \emptyset -definable subsets

Automorphism bases

Definition 0.1 *A subset B of a structure \mathbf{A} is an automorphism base*

\Leftrightarrow

the only automorphism of \mathbf{A} fixing B pointwise is the identity.

Example. $B \subseteq (\mathbb{Q}, <)$ is a-base

\Leftrightarrow

B dense in \mathbb{Q} .

Structures based on \leq_m

- \mathcal{D}_m : all m -degrees
- \mathcal{A}_m : m -degrees of arithmetical sets
- \mathcal{R}_m : c.e. m -degrees

All three structures can be characterized as distributive usl's with a saturation property.

Theorem 0.2 *The minimal degrees do not form an a -base of \mathcal{D}_m , or of \mathcal{A}_m .*

Corollary 0.3 *$\text{Aut}(\mathcal{D}_m)$, $\text{Aut}(\mathcal{A}_m)$ are not simple*

Proof. $\{\pi \in \text{Aut}(\mathbf{A}) : \pi$

fixes the minimal degrees} is a proper normal subgroup. \diamond

A-bases of \mathcal{R}_T

Theorem 0.4 (Ambos-Spies) *For each $c \in \mathcal{R}_T$, $c \neq \mathbf{o}$, $[\mathbf{o}, c]$ is an automorphism base.*

A modified proof of this result [N98] uses the following general method to show that B is an a-base of \mathbf{A} : provide a definable (relative to B)

$$H : \mathbf{A} \mapsto \tau B,$$

where τB is the collection of objects of type τ constructed from B .

τ -maps

Let $\hat{x} := [\mathbf{o}, \mathbf{x}]$. *Example of such a map:*

- $B = \hat{c}$
- $H(\mathbf{x}) = \{\hat{y} \cap \hat{c}, \hat{z} \cap \hat{c} : \mathbf{y} \vee \mathbf{z} = \mathbf{x}\}$.

In general:

- For each formula $\varphi(\mathbf{x})$, we have

$$H_\varphi(\mathbf{x}) = \{\mathbf{x} \in B : \varphi(\mathbf{x})\}$$

- If we have H_1, \dots, H_n , for each formula $\varphi(\mathbf{x}, \mathbf{y}_1, \dots, \mathbf{y}_n)$ have also $H(\mathbf{x}) = \{(H_1(\mathbf{y}_1), \dots, H_n(\mathbf{y}_n)) : \varphi(\mathbf{x}, \mathbf{y}_1, \dots, \mathbf{y}_n)\}$.

A-bases of \mathcal{R}_m

Theorem 0.5 *Each definable $D \subseteq \mathcal{R}_m$, $D \not\subseteq \{\mathbf{0}, \mathbf{1}\}$ forms an a-base of \mathcal{R}_m .*

Proof. Do it only for the set Min of minimal degrees in \mathcal{R}_m . Let

$$\gamma(\mathbf{e}) \equiv \forall \mathbf{q}[\mathbf{e} \not\leq \mathbf{q} \Rightarrow \exists \mathbf{m} \in Min (\mathbf{m} \leq \mathbf{e} \ \& \ \mathbf{m} \not\leq \mathbf{q})]$$

If $\hat{e} \cong \mathcal{R}_m$ effectively, then $\gamma(\mathbf{e})$. Let

$$H(\mathbf{x}) = \{[\mathbf{0}, \mathbf{y}] \cap Min : \mathbf{y} \geq \mathbf{x} \ \& \ \gamma(\mathbf{y})\}$$

- If $z \not\leq x$, then using Denisov's characterization, there is $y \geq x, y \not\leq z$ such that $\gamma(y)$.
- Then, if $H(z) = H(x)$, there must be an $u \geq z, \gamma(u)$, s.t. $[\mathbf{o}, u] \cap Min = [\mathbf{o}, y] \cap Min$, contradiction.

Question 1 *What is the structure of $Aut(\mathcal{R}_m)$?
Is it a simple group?*

From $\mathcal{R}_m - \{\mathbf{1}\} \cong$ the Δ_2^0 m -degrees (Denisov), we can infer the existence of an automorphism of order 2.



The property $\gamma(\mathbf{e})$ plays a role similar to Ambos-Spies' "downward splitting property": let $DSP(\mathbf{a}) \Leftrightarrow$

$$\forall b \not\leq \mathbf{a} \forall c \not\leq \mathbf{b} \vee \mathbf{a}[\hat{\mathbf{c}} \cap \hat{\mathbf{a}} \neq \hat{\mathbf{c}} \cap \hat{\mathbf{b}}].$$

To prove $\hat{\mathbf{c}}$ is an a-base, use as $H(\mathbf{x})$ a complex object built up from ideals $\hat{\mathbf{c}} \cap \hat{\mathbf{a}}$, where $DSP(\mathbf{a})$: for $\mathbf{x} \in NC$, let $H(\mathbf{x}) =$

$$\{ \{ \langle \hat{\mathbf{c}} \cap \hat{\mathbf{a}} : \mathbf{a} \in \hat{\mathbf{p}}_0 \cap DSP \rangle, \langle \hat{\mathbf{c}} \cap \hat{\mathbf{a}} : \mathbf{a} \in \hat{\mathbf{p}}_1 \cap DSP \rangle \} :$$

$$\mathbf{p}_0 \vee \mathbf{p}_1 = \mathbf{p}, \mathbf{p}_i \in NC \},$$

$$\{ \{ \langle \hat{\mathbf{c}} \cap \hat{\mathbf{a}} : \mathbf{a} \in \hat{\mathbf{q}}_0 \cap DSP \rangle, \langle \hat{\mathbf{c}} \cap \hat{\mathbf{a}} : \mathbf{a} \in \hat{\mathbf{q}}_1 \cap DSP \rangle \} :$$

$$\mathbf{q}_0 \vee \mathbf{q}_1 = \mathbf{q}, \mathbf{q}_i \in NC \} :$$

$$\mathbf{p} \vee \mathbf{q} = \mathbf{x}, \mathbf{p}, \mathbf{q} \in NC \}$$

A-bases of \mathcal{E}^*

Theorem 0.6 $Max = \{M^* : M$
maximal $\}$ is an a-base.

Proof. $H(x) =$

$\{[y_0, 1] \cap Max^*, [y_1, 1] \cap Max^* : y_0 \vee y_1 = x\}$. \diamond

Since each maximal element of \mathcal{E}^* is determined by its recursive subsets, it follows that Rec^* is an a-base.

Also the creative sets are (each infinite c.e. set is the union of two creative ones).

Question 2 *Is each orbit of \mathcal{E}^* an a-base ?*

Part II: BI with \mathbb{N} in params

Definition 0.7 *A copy of \mathbb{N} coded in \mathbf{A} with params is a structure*

$$\mathbf{M} = (D, R_+, R_\times) \cong \mathbb{N}$$

such that D and the relations R_+, R_\times are param-definable in \mathbf{A} .

Use the letter \mathbf{M} for such coded structures.

Definition 0.8 *\mathbf{A} is biinterpretable with \mathbb{N} in params*

\Leftrightarrow

$\exists \mathbf{M} \exists f : \mathbf{A} \mapsto \mathbf{M} [f \text{ 1-1, param-def}]$.

Known for $\mathcal{D}(\leq \emptyset')$ (Slaman/Woodin).

Consequences

- Can define a copy $N(\mathbf{A})$ of \mathbb{N} without params
- \mathbf{A} is a prime model and
and a minimal model of $\text{Th}(\mathbf{A})$. In fact, for
 $\mathbf{B}, \mathbf{C} \equiv \mathbf{A}$,

$$\mathbf{B} \prec \mathbf{C} \Leftrightarrow N(\mathbf{B}) \prec N(\mathbf{C}).$$

- Arithmetical \Leftrightarrow param-definable
- Each $\pi \in \text{Aut}(\mathbf{A})$ is arithmetical. Finite
a-base.

(Recall: \mathbf{A} prime \Leftrightarrow each n -orbit
definable, and \mathbf{A} minimal \Leftrightarrow there
is no proper elementary submodel.)

$$\mathcal{R}_m$$

BI fails for \mathcal{R}_m . In fact,

- $|\text{Aut}(\mathcal{R}_m)| = 2^\omega$
- $\exists e[\mathbf{0}, e) \prec \mathcal{R}_m - \{\mathbf{1}\}$

Question 3 *Is \mathcal{R}_m prime ?*

Theorem 0.9 *\mathcal{R}_m prime $\Rightarrow \exists k$ [each realized type is principal via a Σ_k^0 -formula]*

Theorem 0.10 *The arithmetical m -degrees form a prime model.*

The BI-conjecture for \mathcal{R}_T

Question 4 *Is \mathcal{R}_T biinterpretable with \mathbb{N} in params ?*

Weaker:

Question 5 *Is \mathcal{R}_T prime ? ω -homogeneous ?
Minimal ?*

An approx to BI without params

Theorem 0.11 (with Shore, Slaman) *There is a \emptyset -definable*

$$f : \mathcal{R}_T \mapsto N(\mathcal{R}_T)$$

such that

$$\mathbf{x}^{(2)} \neq \mathbf{y}^{(2)} \Rightarrow f(\mathbf{x}) \neq f(\mathbf{y}).$$

Using a-bases

Suppose $B \subseteq \mathcal{R}_T$ is param-definable,
 $H : \mathcal{R}_T \mapsto \tau B$ is 1-1, definable. Then to show BI
it suffices to produce M and a 1-1
param-definable $g : B \mapsto M$:

$$\mathcal{R}_T \xrightarrow{H} \tau B \xrightarrow{\tau g} \tau M \mapsto M.$$

Even if no H is known, we still obtain a finite
a-base: the params for M, g .

Separating sequences

Definition 0.12 A sequence $(g_i)_{i \in \omega}$ separates B from

- *below* $\Leftrightarrow \forall \mathbf{x}, \mathbf{y} \in B$
 $\mathbf{x} \not\leq \mathbf{y} \Rightarrow \exists i (g_i \leq \mathbf{x}, g_i \not\leq \mathbf{y}),$
- *from above* $\Leftrightarrow \forall \mathbf{x}, \mathbf{y} \in B$
 $\mathbf{x} \not\leq \mathbf{y} \Rightarrow \exists i (\mathbf{x} \not\leq g_i, \mathbf{y} \leq g_i).$

M separates B if the sequences $(i^M)_{i \in \omega}$ does. In this case, obtain a 1-1 param-definable $g : B \mapsto M$: e.g. for downward sep. M ,
 $g(\mathbf{x}) = \text{arithm. index for}$

$$\{i : i^M \leq \mathbf{x}\}.$$

Final segments

Theorem 0.13 (Shore, Slaman) • *There is a finite injury construction for M , $\mathbf{b} < \mathbf{1}$ s.t. M separates $[\mathbf{b}, \mathbf{1}]$.*

- *M exists for each promptly simple degree \mathbf{b} .*

It follows that the finite sets of ps degrees are unif'ly definable. The domain of M is a SW-set:

$M = \{x \leq c \text{ minimal} :$

$q \leq x \vee p\}$.

Question 6 *Is there a nontrivial final segment of \mathcal{R}_T which forms an a-base ?*

Question 7 *Is there a u.c.e. antichain which forms an a-base ?*

Definable antichains

Other ways to define antichains:

- $\{x \leq c \text{ maximal} : q \not\leq x \vee p\}$
(Harrington, Shelah; Harrington, Slaman)
- maximum \mathfrak{a} -cappable; only known to yield arbitrarily large finite \mathfrak{a} -chains (A-Sp, Shore, Hirschfeld) but works in all intervals
- $\{x \leq c, d \text{ max} : x \text{ not top of a diamond}\}$
(A-Sp, Soare; Lempp, N)

Initial segments

- Have to separate upward
- Problem: the constructions yield u.c.e. antichains.

Proposition 0.14 *If (g_i) is a u.c.e. sequence of nonzero degrees and $c \neq \mathbf{o}$, then*

$$\exists v < u \leq c \forall i [u \vee g_i = v \vee g_i].$$

Possible solutions:

- Find a more flexible construction of a definable a-chain
- use different types of interaction of M and $[\mathbf{o}, c]$
- forget it.

Approximations

Theorem 0.15 (N, ?) *There is*

- $c \neq \mathbf{0}$
- $X \subseteq \hat{c}$
- $f : X \mapsto M$ 1-1

such that

$$\forall u \leq c, u \neq \mathbf{0} \exists x \in X$$

$$[x \neq \mathbf{0} \ \& \ x \leq u]$$

Proof.

- Construct c.e. sets C, D
- Let $X =$ degrees of c.e. set splits of C
- Construct M separating X from above.
Domain of M is a SW set.

The advantage of set splits U : control size of U changes in reaction to C changes.

If, in addition, we make C nonbounding, we obtain a strictly descending definable sequence (\mathbf{b}_i) s.t.

$$\forall \mathbf{x} \leq \mathbf{c}[\mathbf{x} \neq \mathbf{o} \Rightarrow \exists i \mathbf{b}_i \leq \mathbf{x}].$$

Part III: \emptyset -definability

Examples of interesting \emptyset -definable subsets of \mathcal{R}_T :

- Promptly simple degrees (A-Spies e.a.)
- L_n, H_{n-1} for $n \geq 2$ (N, Shore, Slaman)
- Contiguous degrees (Downey, Lempp)

Question 8 *Is there a \emptyset -definable ideal except the cappable degrees and the noncappable degrees ?*

Weaker, but still interesting: proper ideals generated by a definable set.

Theorem 0.16 *The ideal generated by NB is properly contained in CAP.*

Local structure of d'ble sets

Question 9 (Li-Ansheng) *Is there a \emptyset -definable D and an interval $\mathbf{a}, \mathbf{b}]$ such that $\mathbf{a}, \mathbf{b} \notin D$ and $|\mathbf{a}, \mathbf{b}] \cap D| = 1$?*

Good candidates for D : contiguous; sup of a minimal pair. *Note: probably solved by Cholak, Downey and Walk*

\emptyset -d'ty of a copy of \mathbb{N}

Can we define a copy $N(\mathbf{A})$ of $(\mathbb{N}, +, \times)$ in \mathbf{A} without params ?

- Possible for \mathcal{R}_m [N95], \mathcal{R}_T [N,Shore, Slaman 96] and for \mathcal{R}_{wtt} [Nta].
- Fails for \mathcal{E}^* .
- Unknown for $\mathcal{R}_Q, \mathcal{R}_{tt}$