

# **HASH FUNCTION, MAC, and HMAC**

## **Lecture 5**

**COMPSCI 726**

**Network Defence and Countermeasures**

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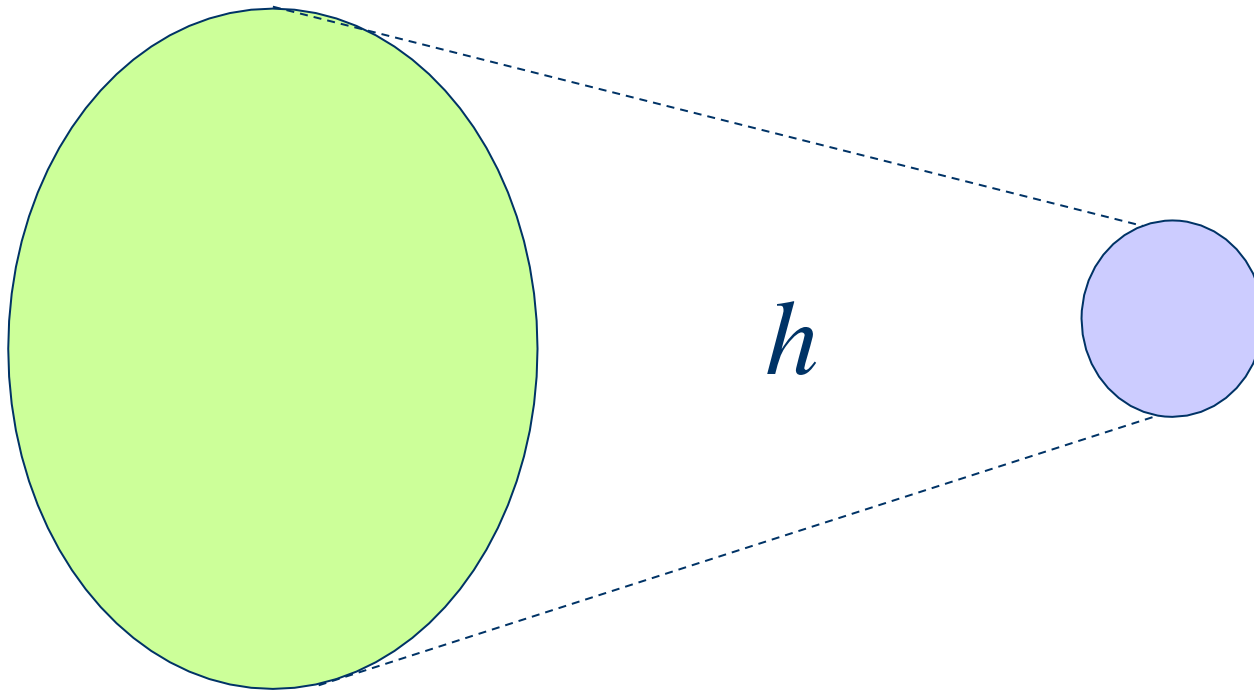
July 28, 2021

Source of some slides: **Stanford University**

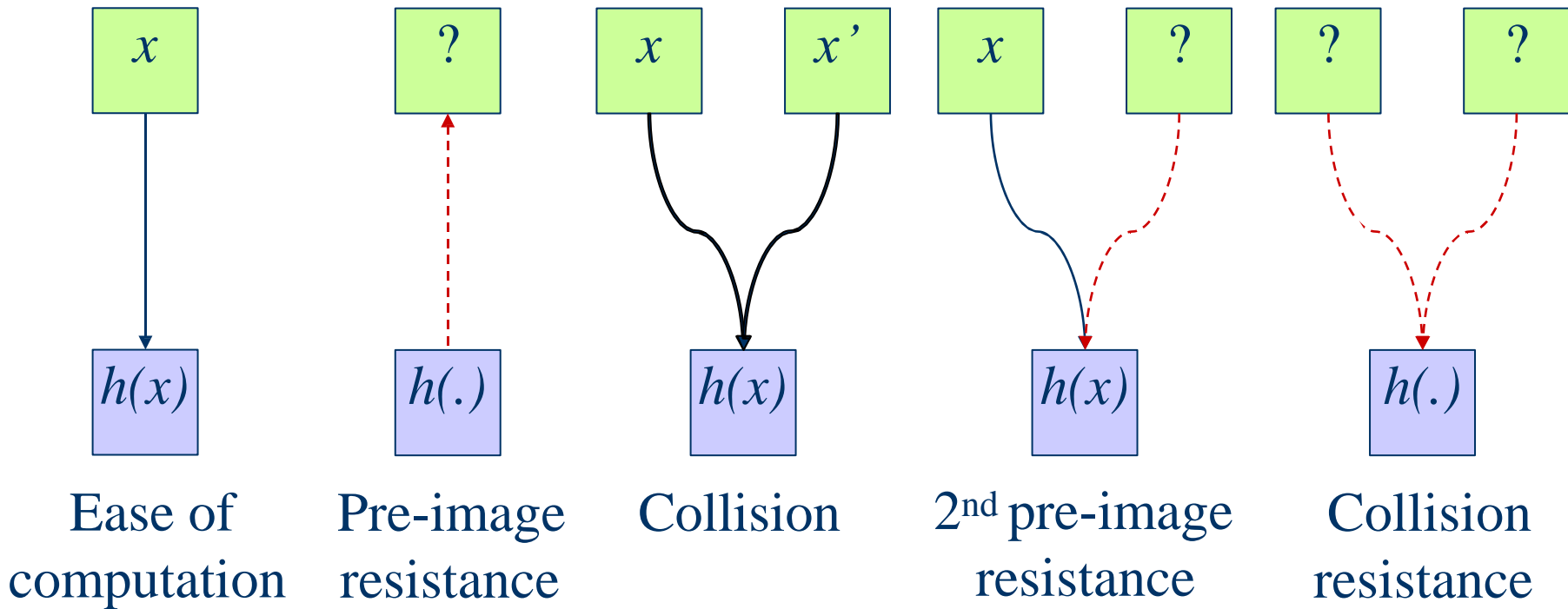


# HASH FUNCTION

- Length-reducing function  $h$ 
  - Maps an arbitrary string to a fixed-length string
- Publicly known
- Also known as cryptographic checksum or message digest

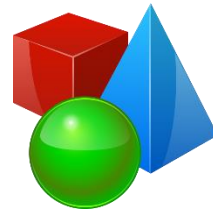


# HASH PROPERTIES



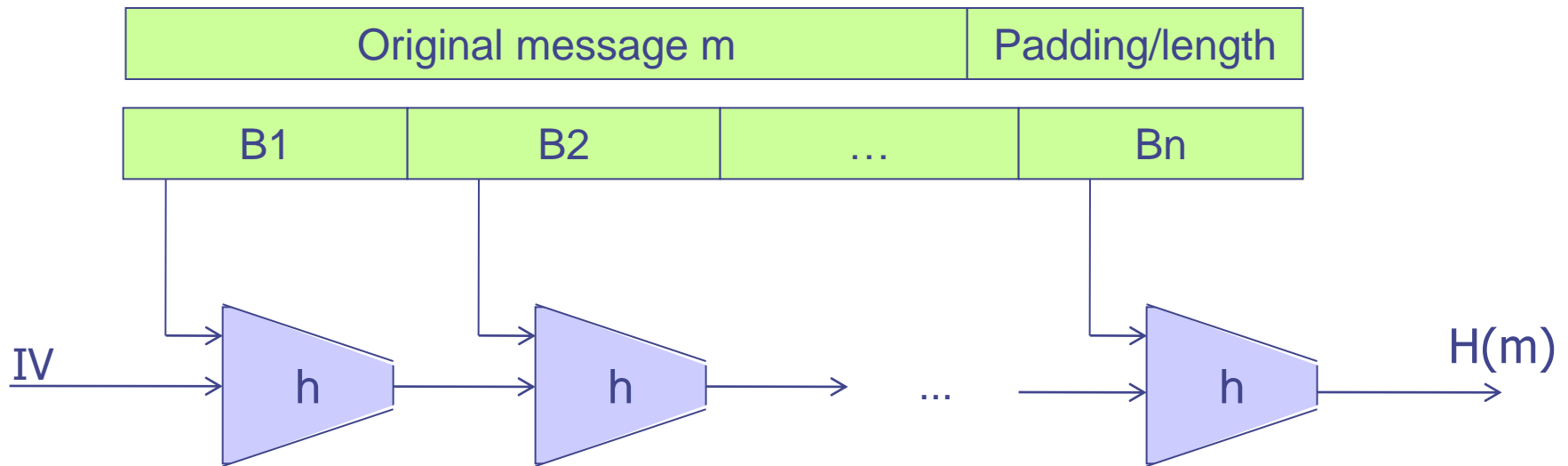
- Collision resistance implies 2<sup>nd</sup> pre-image resistance

# COMMONLY USED HASH FUNCTIONS



- **MD** (Message Digest)
  - MD5
    - Max message  $< 2^{64}$
    - Output: 128-bit
- **SHA** (Secure Hash Algorithm)
  - SHA-1
    - Max message  $< 2^{64}$
    - Output: 160-bit
  - SHA-2
    - Max message  $< 2^{128}$
    - Max output: 512-bit
  - SHA-3
    - Max message: Unlimited
    - Max output: 512-bit

# SHA-512: MERKLE-DAMGARD SCHEME



- Augmented message: multiple of 1024-bit blocks
- $h(., B_i)$  is a compression function
- Theorem: If  $h$  is collision resistant then so is  $H$

# HASH APPLICATIONS



- Detect changes to messages/files (integrity)
- Digital signatures
  - Sign hash of message instead of entire message
- Psudorandom function (PRF)
  - Generate session key, nonce (Number Only Once)
  - Produce key from password
  - Derive keys from master key
- Create one-way password file
  - Store hash of password
    - Salt to harden pre-computed dictionary attacks
- Viruses and intrusion detection
- Auctions: To bid  $B$ , send  $h(B)$  and reveal  $B$  later

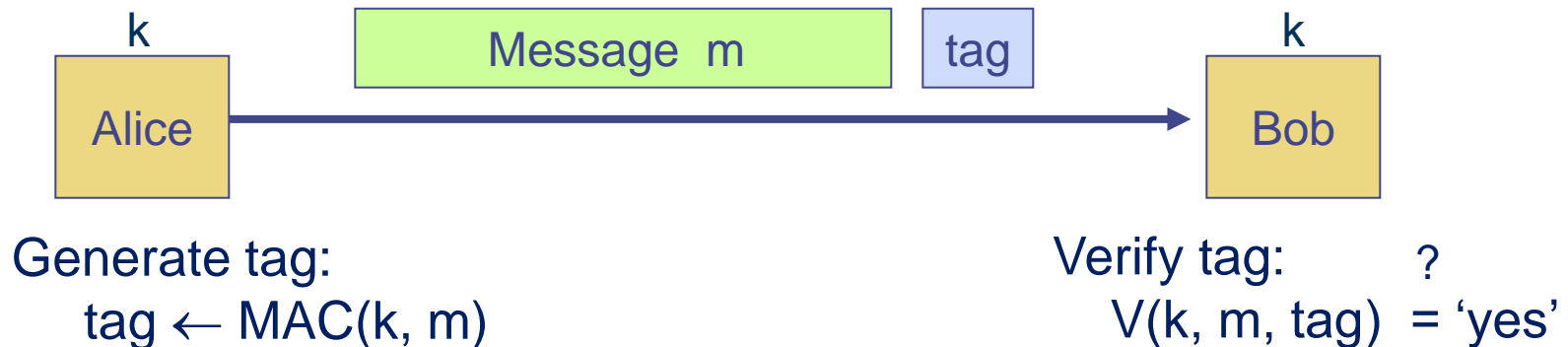
# HASH VS. ENCRYPTION



- Hashing is a one-way
  - No unhashing
- Publicly known and there is no key used
- Efficient
- Deterministic (compared)
  - $H(m) == H(m')$
  - Of course, hashes with salts are not!
    - $H(m \parallel s1)$  and  $H(m \parallel s2)$
- Encryption is not one-way
  - Decryption renders the original message
- Publicly known algorithms but the key is kept secret
- Slower
- May or may not be deterministic (compared)
  - Randomised encryption
    - $Enc(k, t1 \parallel m)$  and  $Enc(k, t2 \parallel m)$

# MESSAGE AUTHENTICATION CODE (MAC)

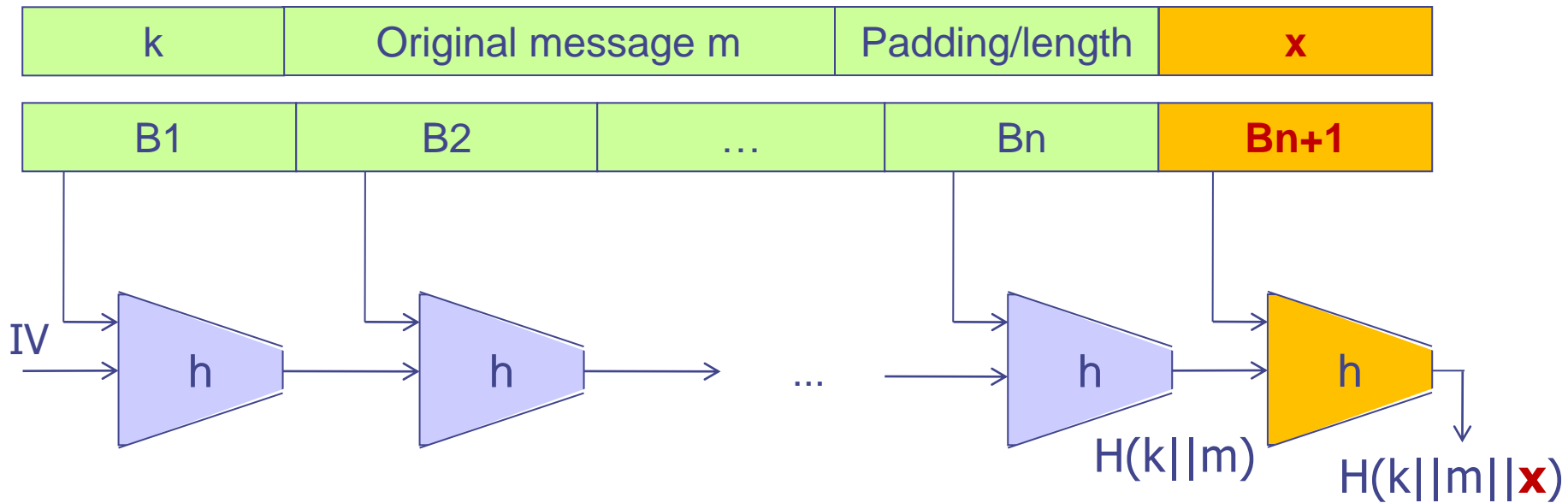
- Like a hash function, but it uses a key!
- Appended to the original message
- Receiver performs same computation on the message and checks if it matches the MAC
- It provides assurance that the message is unaltered and comes from the sender





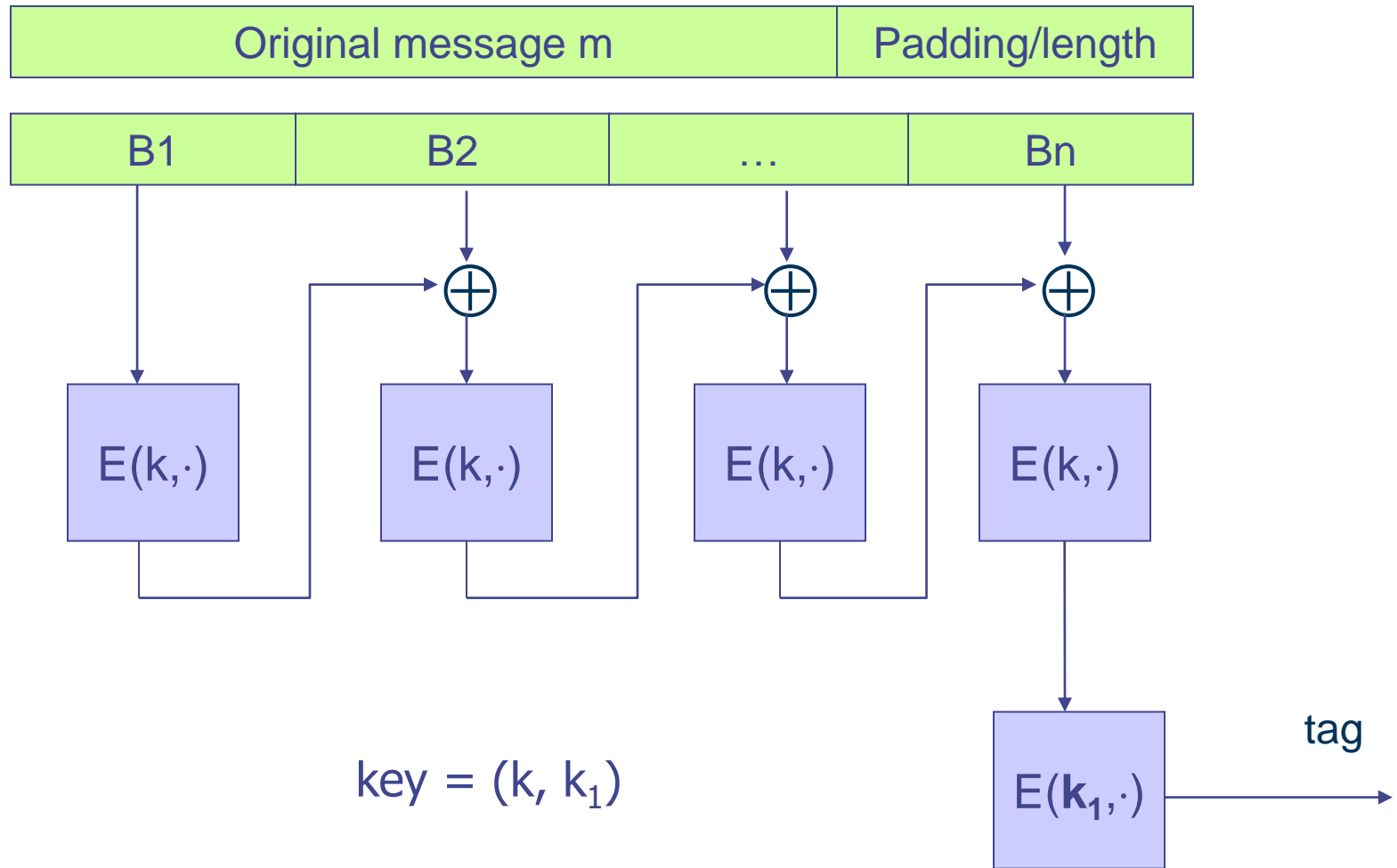
# MAC CONSTRUCTION: MERKLE-DAMGARD SCHEME

- $\text{MAC}(k, m) = H(k \parallel m)$



- Issue: Length extension attack!

# MAC CONSTRUCTION: RAW CBC



# MAC APPLICATIONS



- Integrity of a message or file
- Validating identity of a message sender (authentication)

# HASH VS. MAC



- Publicly known and no key
  - A hash value
  - Efficient
  - Message integrity
  - Anyone can generate it
- Publicly known, but the key is kept secret
  - A keyed hash value
  - Slower
  - Message integrity and authentication
  - Only an authorised user can generate it

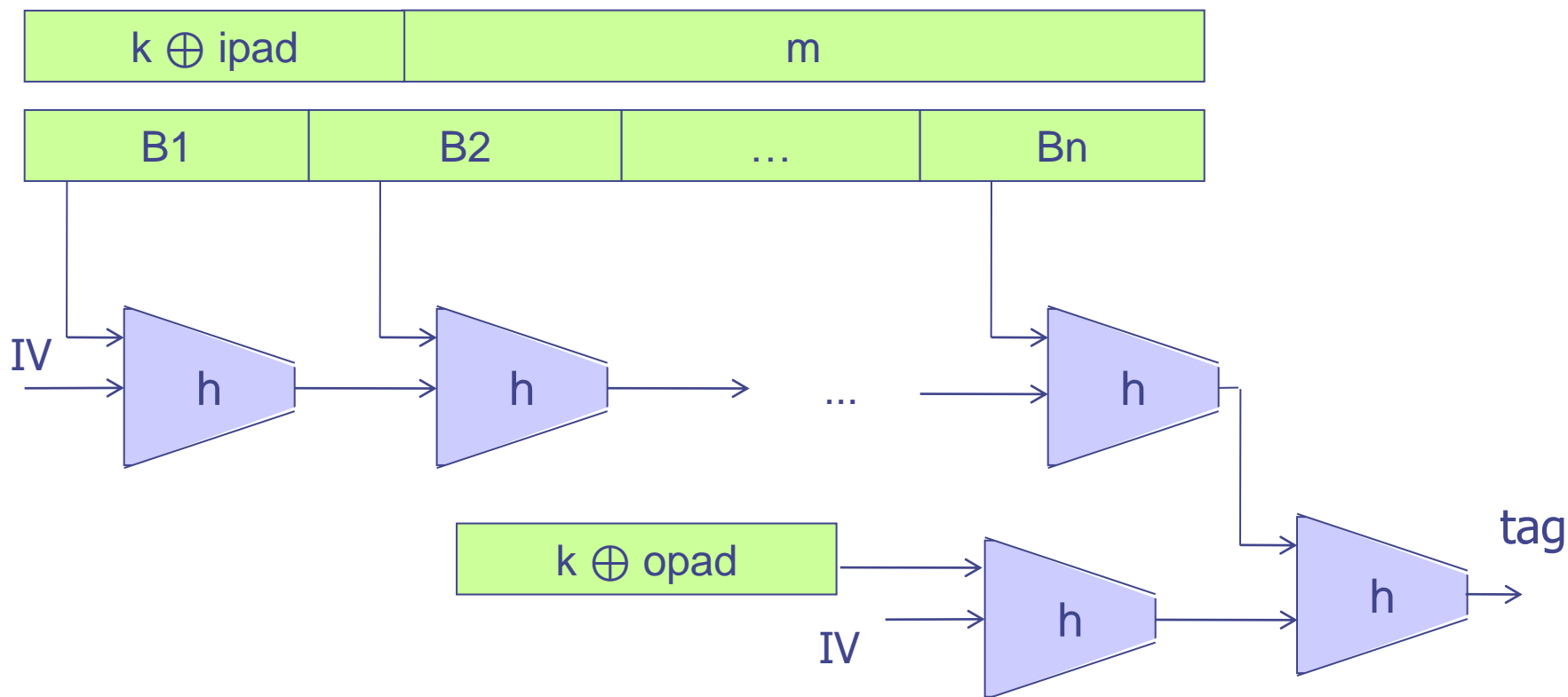
# HASH-BASED MAC (HMAC)



- Evolved from weakness in MAC
- A specific construction of calculating a MAC involving a secret key
- Uses and handles the key in a simple way
- Less effected by collision than underlying hash algorithm
- More secure

# HMAC CONSTRUCTION: MERKLE-DAMGARD SCHEME

- $\text{HMAC}(k, m) = H(k \oplus \text{opad} \parallel H(k \oplus \text{ipad} \parallel m))$



- Theorem: If  $h$  is a PRF then HMAC is a PRF

# AUTHENTICATED ENCRYPTION



Encryption key  $K_E$       MAC key =  $K_I$

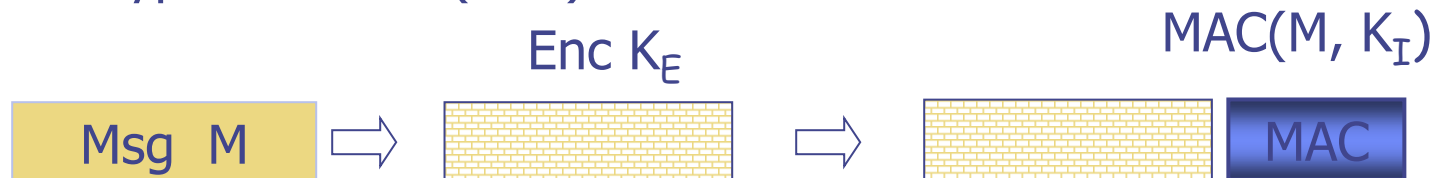
Option 1: MAC-then-Encrypt (SSL)



Option 2: Encrypt-then-MAC (IPsec)



Option 3: Encrypt-and-MAC (SSH)



# SUMMARY



- Hash is a one-way function, which is easy to compute but difficult to invert
- MAC offers both data integrity and authentication
- Authenticated encryption combines both encryption and MAC



# RESOURCES



- Read Chapter 3 of  
**Network Security Essentials – Applications and Standards**  
Fourth Edition  
William Stallings  
Prentice Hall  
ISBN 0-13-706792-5



**Questions?**

**Thanks for your attention!**