IPv4 Lecture 17

COMPSCI 726 Network Defence and Countermeasures

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Slides from Muhammad Rizwan Asghar

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Source of some slides: **Princeton University** Also thanks to **J.F Kurose and K.W. Ross**







- Internet Protocol (IP) is the principal communication protocol in TCP/IP
- IP was proposed in the early 1970s
- IPv4 is IP version 4
- A 32-bit address that uniquely and universally identifies a host on the Internet
 - E.g., 203.118.141.95 (ping <u>www.google.com</u>)

IPV4 ADDRESSING: 5 CLASSES



a. Binary notation

b. Dotted-decimal notation

Source: Data Communications and Networking by Behrouz A. Forouzan

IPV4 ADDRESSING: BLOCKS AND BLOCK SIZE

Class	Number of Blocks	Block Size	Application
А	128	16,777,216	Unicast
В	16,384	65,536	Unicast
С	2,097,152	256	Unicast
D	1	268,435,456	Multicast
Е	1	268,435,456	Reserved

Source: Data Communications and Networking by Behrouz A. Forouzan

IPV4 PACKET FORMAT



IP HEADER FIELDS

 4-bit Version
 4-bit Header
 8-bit Type of Service
 16-bit Total Length (Bytes)

 16-bit Identification
 3-bit Flags
 13-bit 13-bit Fragment Offset

 8-bit Time to Live (TTL)
 8-bit Protocol
 16-bit Header Checksum

 32-bit Source IP Address
 32-bit Destination IP Address

 32-bit Destination IP Address
 Options (if any)

- Version number (4-bit)
 - Indicates the version of the IP protocol
 - Typically 4 (for IPv4) and sometimes 6 (for IPv6)
- Header length (4-bit)
 - Number of 32-bit words in the header
 - Typically 5 (for a 20-byte IPv4 header)
- Type of service (8-bit)
 - Used to manage quality of service
 - E.g., low delay for audio and high bandwidth for bulk transfer

IP HEADER FIELDS CONT



- Total length (16-bit)
 - Number of bytes in the packet (header+payload)
 - Maximum size can be 64KB
 - Underlying links may impose harder limits
- Fragmentation information (32-bit)
 - Packet identification, flags and fragmentation offset (see later)
 - Supports dividing a large IP packet into fragments when a link cannot handle that packet
- Time-to-live (8-bit)
 - Lifetime of a packet
 - Used to prevent loops reduces as the packet traverses

IP HEADER FIELDS CONT



- Protocol (8-bit)
 - A value that specifies the type of payload
 - E.g., TCP or UDP
- Header checksum (32-bit)
 - It includes all other fields in an IP header
 - Recalculated by each router since TTL changes
- Source or destination address (32-bit)
 - IP address



Potential robustness problem

- Forwarding loops can cause packets to cycle forever
- Confusing if the packet arrives later



- TTL in packet header (8-bit)
 - TTL decremented by each router on the path
 - A packet is discarded when TTL reaches 0
 - A 'time exceeded' message is sent to the source

IP SPOOFING



- Source IP address should be the sending host
 - But, who is checking that?
 - You could send packets with any source you want
- Why would someone want to do this?
 - Launch a DoS attack
 - Evade detection
 - An attack against the spoofed host
 - Spoofed host is wrongly blamed
 - Spoofed host may receive return traffic from the receiver

IP FRAGMENTATION AND REASSEMBLY

- Max IP datagram: 64KB
- Network links have Maximum Transfer Unit (MTU)
- Large IP datagrams can be fragmented
- Reassembled at destination



IP FRAGMENTATION AND REASSEMBLY CONT



ID identifies IP datagram

fragflag=1 means more fragments available

offset points fragment offset (in octet)

ISSUES WITH FRAGMENTATION



- Uses resources poorly
 - A bit shorter MTU will require a full fragment and a smaller fragment
- Poor end-to-end performance
- Reassembly is hard
 - Buffering constraints
- Reassembly is slow
- Interferes with TCP control flow



Questions?

Thanks for your attention!