COMPSCI 314 S1 C

Data Communications Fundamentals

Lecture Slides, Set #5

Clark Thomborson 7 April 2006

RFC 2410: "The NULL Encryption Algorithm and Its Use With IPsec"

- "... NULL is a block cipher the origins of which appear to be lost in antiquity.
- "Despite rumors that the National Security Agency suppressed publication of this algorithm, there is no evidence of such action on their part.
- "Rather, recent archaeological evidence suggests that the NULL algorithm was developed in Roman times, as an exportable alternative to Ceaser [sic] ciphers.
- "However, because Roman numerals lack a symbol for zero, written records of the algorithm's development were lost to historians for over two millennia.
- "NULL is defined mathematically by the use of the Identity function I applied to a block of data b such that:

$$NULL(b) = I(b) = b$$

 "Like other modern ciphers, e.g., RC5 [RFC-2040], the NULL encryption algorithm can make use of keys of varying lengths. However, no measurable increase in security is afforded by the use of longer key lengths."

IPsec

- There are two main parts to the IP security protocol:
 - IKE, the Internet Key Exchange, to set up keys
 - ESP/AH, the Encapsulating Security Payload and Authentication Header.
 - (ESP is *not* the "Encryption Security Payload" ;-)
- This is a very complex protocol: "over-engineered"?
 - RFC 2401 gives a 66 pp. overview (obsoleted by RFC 4301, December 2005, 99 pp.).
 - RFC 2402 defines AH; was obsoleted by RFC 4302 (33 pp.) and RFC 4305 (5 pp.)
 - Also: RFC 2403, RFC 2404, RFC 2405, RFC 2410, RFC 2411, RFC 4109, RFC 4306 ...
 - See http://rfc.net/rfc-index.html.

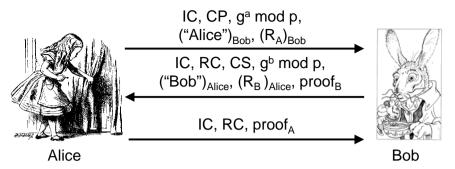
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Internet Key Exchange (IKE)

- There are two phases:
 - Phase 1: set up an IKE "security association". Four options:
 - Public key encryption (original)
 - Public key encryption (improved)
 - Digital signature (for use when Alice doesn't initially know Bob's public key)
 - Symmetric Key
 - For each of these options there is a "main mode" (6 messages) and an "aggressive mode" (3 messages, no anonymity for Bob or Alice except in public key options). Eight versions!!
 - Phase 2: set up an IPsec "security association"
 - This is a 3-message key-exchange.
- Source: M Stamp, Information Security: Principles and Practice, Wiley, 2006.

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IKE PK, Aggressive Mode



$$\begin{split} & \text{IC = initiator cookie,} & \text{CP = crypto proposed,} \\ & \text{CS = crypto selected,} & \text{RC = responder cookie,} \\ & \text{SKEYID = h(R_A, R_B, g^{ab} \bmod p),} \\ & \text{proof}_{\text{A}} = [\text{h(SKEYID, g}^{a} \bmod b, g^{b} \bmod p, \text{IC, RC, CP,} \\ & \text{"Alice"})]_{\text{Alice}} \\ & \text{Source: M Stamp, Informatical Proofs and Position of the proofs and Proofs a$$

Source: M Stamp, Information Security Principles and Practice, Wiley, 2006.

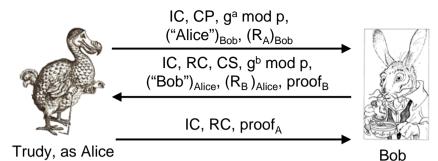
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IPsec cookies

- The IC and RC cookies are called "anti-clogging tokens" in the relevant RFCs.
- These are unrelated to web cookies, which maintain state in a web browser.
- The IC and RC cookies are intended to allow Bob to remain stateless as long as possible during a session initiation.
 - In a typical DoS (denial-of-service) attack, Bob is overwhelmed by a flood of spurious requests for session initiations.
 - Bob will run out of memory if he tries to remember all session initiations that haven't yet completed.
- Bob must remember CP from the first IKE message, when responding to the third message.
 - Bob is susceptible to DoS attacks on IKE initiations.

Source: M Stamp, Information Securit Principles and Practice, Wiley, 2006.

Aggressive IKE PK is anonymous (!)



$$\begin{split} & \text{IC} = \text{initiator cookie}, & \text{CP} = \text{crypto proposed}, \\ & \text{CS} = \text{crypto selected}, & \text{RC} = \text{responder cookie}, \\ & \text{SKEYID} = h(R_A, R_B, g^{ab} \text{ mod p}), \\ & \text{proof}_A = [h(\text{SKEYID}, g^a \text{ mod b}, g^b \text{ mod p}, \text{IC}, \text{RC}, \text{CP}, \text{``Alice''})]_{Alice} \end{split}$$

- This is called "plausible deniability" in IKE.
- · Alice can deny sending messages to Bob: "Trudy did it!"

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IPsec datagrams

- An IPsec datagram is protected by one of two protocols.
 - ESP (the Encapsulating Security Payload) protects integrity and confidentiality.
 - AH (the Authentication Header) protects integrity.
 - Everything beyond the IP header is protected.
- IPsec defines two transport modes.
 - In transport mode, the new ESP/AH header is sandwiched between the IP header and the data.
 - In tunnel mode, the entire IP packet is encapsulated in a new IP packet.

Source: M Stamp, Information Security: Principles and Practice, Wiley, 2006.

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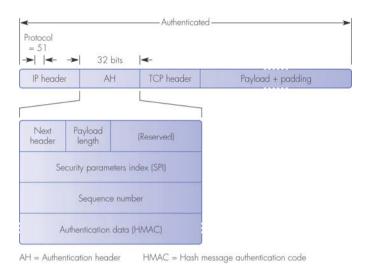


Figure 10.15 AH protocol header position and contents in transport mode

What can an attacker learn by analysing this packet?

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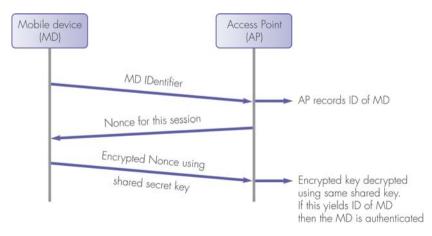
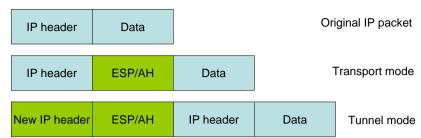


Figure 10.17 Authentication of a mobile device by the AP using a three-way handshake procedure

Has the AP authenticated itself to the MD?

Transport Mode vs. Tunnel Mode



- Transport mode adds less header information:
- Source: M Stamp, Information Security. Principles and Practice, Wiley, 2006.

- more efficient.
- Tunnel mode with ESP doesn't let the attacker see the IP header.
 - Firewall-to-firewall packets don't reveal intranet addresses.
 - The attacker can see firewall (internet) addresses in the new IP header.

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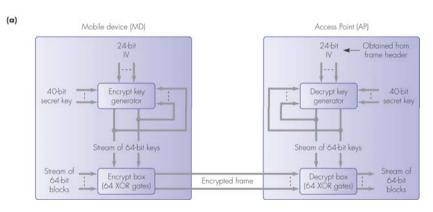


Figure 10.18 WEP protocol: (a) RC4 stream cipher principles

Note that the IV is sent "in the clear"! Eavesdropper can see when the same IV is used twice => same stream cipher => enables a cryptographic attack on RC4.

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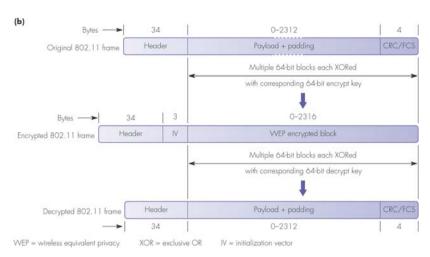


Figure 10.18 WEP protocol: (b) frame encryption/decryption operations

Note use of CRC: a poor choice with XOR (stream cipher), allows modification attacks.

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802.11g, 802.11i, ...

- 802.11g: defines Wi-Fi Protected Access (WPA).
 - Ratified as a standard in June 2003.
 - Includes some of the security improvements in a draft of 802.11i.
- 802.11i: defines WPA2.
 - Ratified in June 2004.
 - Uses the Advanced Encryption Standard (AES), instead of RC4, which was used in WEP and WPA.

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- 802.11w: protect management and broadcast frames.
 - Task group started work in January 2005.
 - Target date for ratification: March 2008.
- Sources:

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- http://en.wikipedia.org/wiki/802.11 (23:30, 4 April 2006),
- http://en.wikipedia.org/wiki/IEEE_802.11w (07:14, 7 March 2006).

In(Security) of the WEP Algorithm

- A busy access point, which constantly sends 1500 byte packets at 11Mbps, will exhaust the space of IVs after $1500*8/(11*10^6)*2^24 = \sim 18000$ seconds, or 5 hours.
 - The amount of time may be even smaller, since many packets are smaller than 1500 bytes.
- This allows an attacker to collect two ciphertexts that are encrypted with the same key stream and perform statistical attacks to recover the plaintext.
- Worse, when the same key is used by all mobile stations, there are even more chances of IV collision.
 - For example, a common wireless card from Lucent resets the IV to 0 each time a card is initialized, and increments the IV by 1 with each packet.
 - This means that two cards inserted at roughly the same time will provide an abundance of IV collisions for an attacker.
- Source: http://www.isaac.cs.berkeley.edu/isaac/wep-faq.html, 2001.

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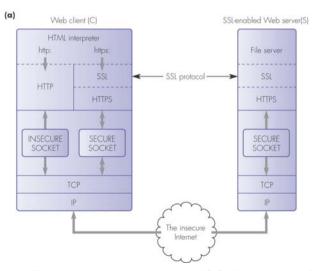


Figure 10.19 The secure socket layer (SSL) protocol: (a) protocol stack

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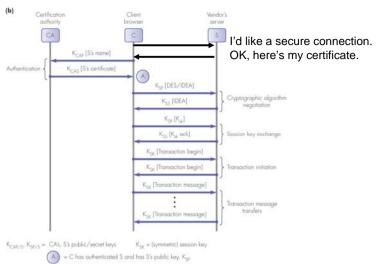


Figure 10.19 The secure socket layer (SSL) protocol: (b) outline of the authentication and transaction initiation phases

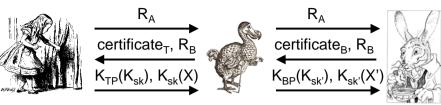
Does C authenticate to S? Does S authenticate to C? What if CA is offline?

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Secure Electronic Transfer (SET)

- Developed by Visa & Mastercard, starting in 1996.
- Excellent cryptographic design, based on iKP family of protocols (i=1, 2, 3) from IBM R&D Zurich. Assurances:
 - Confidentiality (merchant doesn't learn customer's credit card number); Integrity;
 - Authentication of merchants (for i>1), and of customers (i=2).
 - Non-repudiability by merchants (for i>1), and by customers (i=2).
 - Source: Bellare et al., *IEEE J SAC 18:4* 611-627, April 2000.
- Failed to win market share. Visa now uses 3-D Secure.
 - For i=2, customers must have a public key certificate, e.g. on a smart-card. (A chicken-and-egg problem.)
 - Doesn't support e-commerce from home (smart-card reader expense).
 - Source: http://cs.bilkent.edu.tr/~selcuk/teaching/cs519/cs519.21.ppt
 - Some merchants use credit card numbers to recognize repeat customers. (A legacy-software problem.)

SSL Protection against Man-in-the-Middle



Alice

Trudy: acting as Alice to Bob, and as Bob to Alice

3ob

- How can Alice detect that Trudy is "in the middle"?
- What does your web-browser do, when it receives an SSL certificate that is invalid?
 - Your browser has a collection of public keys for CAs.
- Have you ever inspected an SSL certificate?

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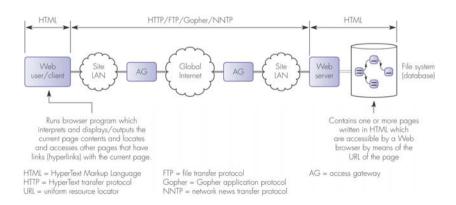


Figure 9.1 Information browsing

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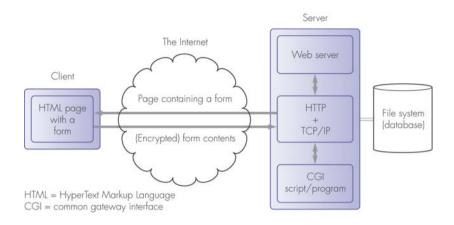


Figure 9.3 Electronic commerce

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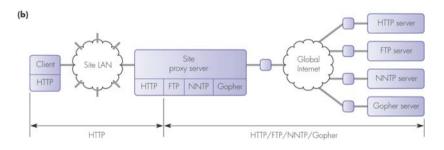


Figure 9.4 Intermediate systems: (b) proxy server

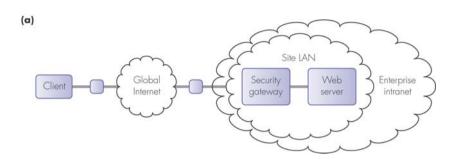


Figure 9.4 Intermediate systems: (a) security gateway

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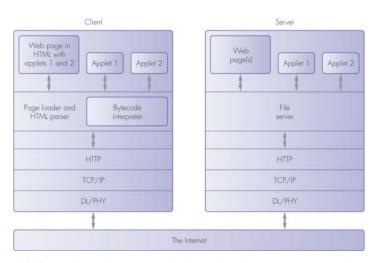


Figure 9.5 Protocol stack to support the browsing of pages containing Java applets

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