

CS314s2 - lectures 22-31

How the Internet works

Dr Brian Carpenter

- 22, 23 – Routing
 - 24, 25 – IP version 4, DNS
 - 26 – IP version 6
 - 27 – Transport protocols: TCP
 - 28 – Transport protocols: UDP; Sockets
 - 29 – Applications (SSH, FTP, SMTP,...)
 - 30 – Applications (HTTP and the Web)
 - 31 – Summary
-
- Dates: 25, 27, 28 Sept., 2, 4, 5, 9, 11, 12, 16 Oct.
 - Approximately covering Shay chapters 10.4 to 12
 - **Questions:** brian@cs.auckland.ac.nz or room 303s.587 (most days between 10 a.m. and 4 p.m.)

About the level of detail

- I tend to give more technical details than the text books, because I want all students to feel that they can see how the various protocols could be coded in a programming language.
 - I will also give some advanced references, which are for interest only and quite optional.
- The level of detail given in Shay or Halsall should be sufficient for the exam questions.

CS314s2- 22/23

Routing algorithms and protocols

- This is quite a hard topic.
- In two lectures, we can only scratch the surface to explain principles.

Topics

- Routing – the problem
- Approaches
- Algorithms
 - Dijkstra
 - Bellman-Ford
- Protocols
 - RIP
 - OSPF
 - BGP
- Final points

Routing – related to distance

- How do we connect devices

- In the same room?

- In the next room? (tens)

- In the next building? (100s)

- In the next department?

- In the next campus? (1000s)

- In the next country? (millions)

- In the next continent? (billions)

Cable/radio boundary

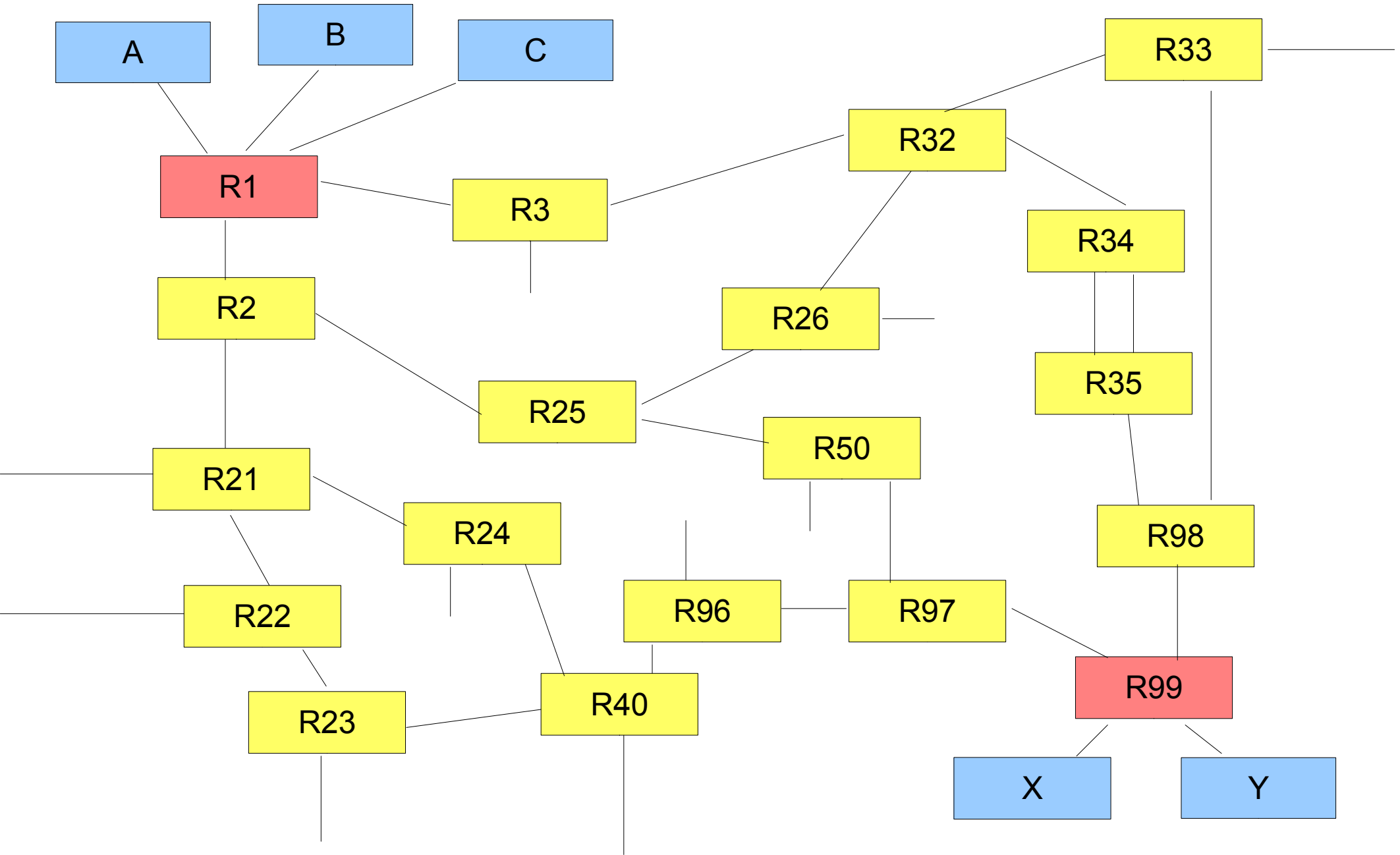
+ Admin boundary

+ Funding boundary

+ Political boundary

Distance, logical
separation, cost and
numbers increase; must
share costs and cables

Routing – the problem (1)



Terminology note

- This sort of picture shows what mathematicians call a *graph*, and there is a lot of *graph theory*.
- A graph consists of numbered *vertices* connected by *arcs*:

$V_1 \rightarrow V_2$ would be the arc from vertex 1 to 2

- For a program, a graph can be thought of as an array of vertices and a matrix of arcs
 - of course, not all arcs exist; $R2 \rightarrow R3$ does not exist so its matrix entry will be zero.
- We will not dig deeply into graph theory here

Graphs in programs

- Representing a graph:
 - Ordered list of vertex names:
 $V_1 = 'A', V_2 = 'R1', \dots, V_{22} = 'R99', V_{23} = 'X', V_{24} = 'Y'$
 - List of arcs (i.e. a sparse matrix)
 $A_i = (V_s, V_d, W)$
 - (source, destination, and a parameter for each arc)
- A path through the graph is represented by an ordered list of arcs
 - “Walking the graph” means an algorithm that starts at a given vertex and builds paths by searching the list of arcs that meet up (the destination of one arc is the source of the next).

Routing – the problem (2)

- Choose the best path from A to X, knowing (at A) only the logical address of X.
- “Best” could mean
 - Smallest number of hops
 - Shortest time delay
 - Least congested
 - Cheapest
 - Administratively allowed
 - Easiest to discover
 - Any combination of the above
- Solution must be reasonably quick and guaranteed to avoid loops and deadlocks

Routing – the problem (3)

- The real network is massive
 - Thousands of nodes on campus, millions per country, hundreds of millions worldwide.
- The real network is constantly changing
 - Nodes, routers and links added or removed constantly
 - View of network will be different in different places
- Therefore, algorithms must be scaleable, dynamic and distributed
 - Central static route computation only works for small, very stable networks

Approaches to dynamic routing

- Confine routing intelligence to specialised router boxes
 - End-systems mainly send everything to their default router
- Routers build connectivity maps and exchange routing information with their neighbours
- Two major approaches
 - Distance vector routing: router informs neighbours of its best paths. Each router calculates best paths for itself based on neighbour's paths.
 - Link state routing: router exchanges link status info with its neighbours. Each router relies on accuracy of its neighbours' status info in calculating best paths through the whole map.

DV Routing: Bellman-Ford algorithm

- Router computes a table of its “distance” to each other router in the same network
 - Distance 1 = 1 hop, 2 = 2 hops etc.
- Sends this table to its neighbours
- Recomputes its own table using tables from its neighbours
- Works, but
 - Sending complete tables doesn't scale well and is slow to converge after a change
 - “Link down” condition can lead to loops, as the “distance” increases towards infinity

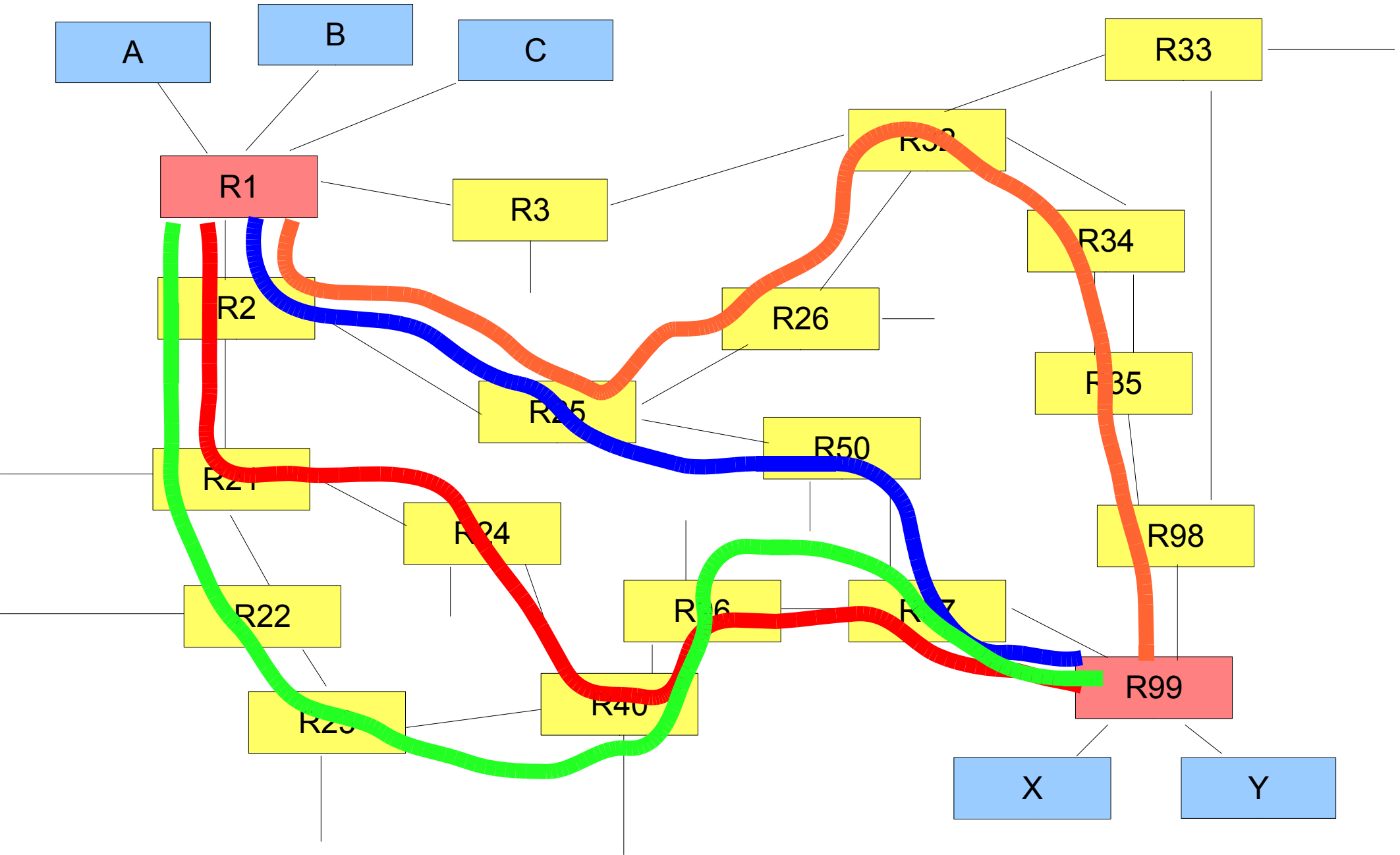
Bellman-Ford: the computation (1)

- Compute the shortest path from R_1 to all R_n knowing the topology of the router graph
 - Initialize each distance $D(R_1, R_n) = \infty$, except $D(R_1, R_1) = 0$
 - For each arc $R_x \rightarrow R_y$ in the graph (starting with R_1)
 if $D(R_1, R_y) > D(R_1, R_x) + 1$
 then $D(R_1, R_y) := D(R_1, R_x) + 1$
 - At this point, each $D(R_1, R_n)$ is the length of the shortest path from R_1 to R_n
 - The algorithm also keeps track of the “next hop” router from R_1 for the path

Bellman-Ford: the computation (2)

- That dealt with paths from R1. For complete routing tables, run again starting from R2, R3...
- Note that algorithm processes every arc times every node!
- The “+1” above is a simplification where all paths have equal weight. It should be a weight $W(x, y)$ to allow for administrative preferences for certain links:
$$\text{if } D(R1, Ry) > D(R1, Rx) + W(x, y)$$
$$\text{then } D(R1, Ry) := D(R1, Rx) + W(x, y)$$
- $W(x, y) < 0$ is allowed by the algorithm, but can lead to loops, which can be detected if any $D(R1, Ry) > D(R1, Rx) + W(x, y)$ condition exists after the algorithm ends

Bellman-Ford paths



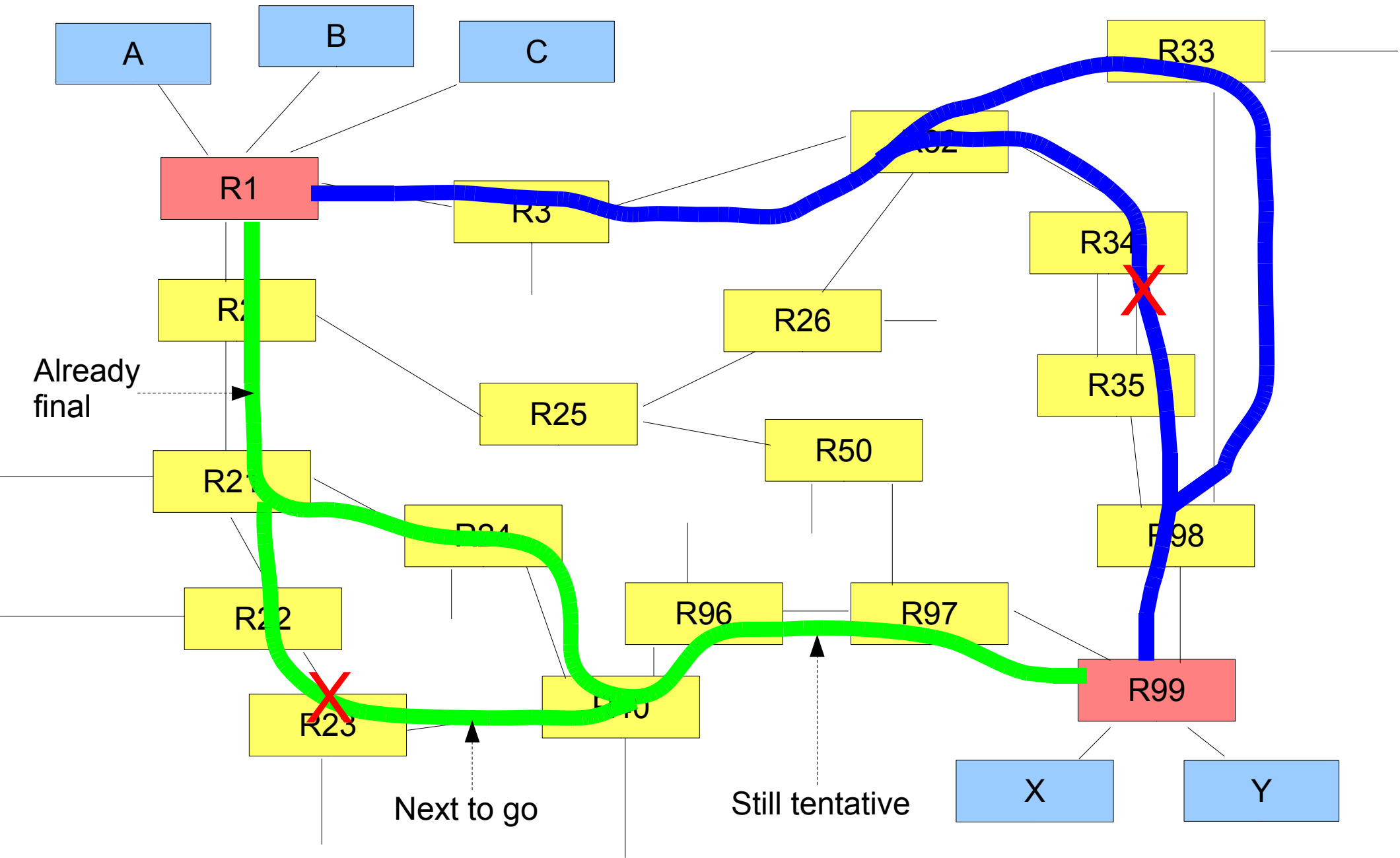
Link State Routing: Dijkstra algorithm

- Also called “shortest path first” or “SPF”
- First we'll describe the algorithm in the abstract
- Then a few comments

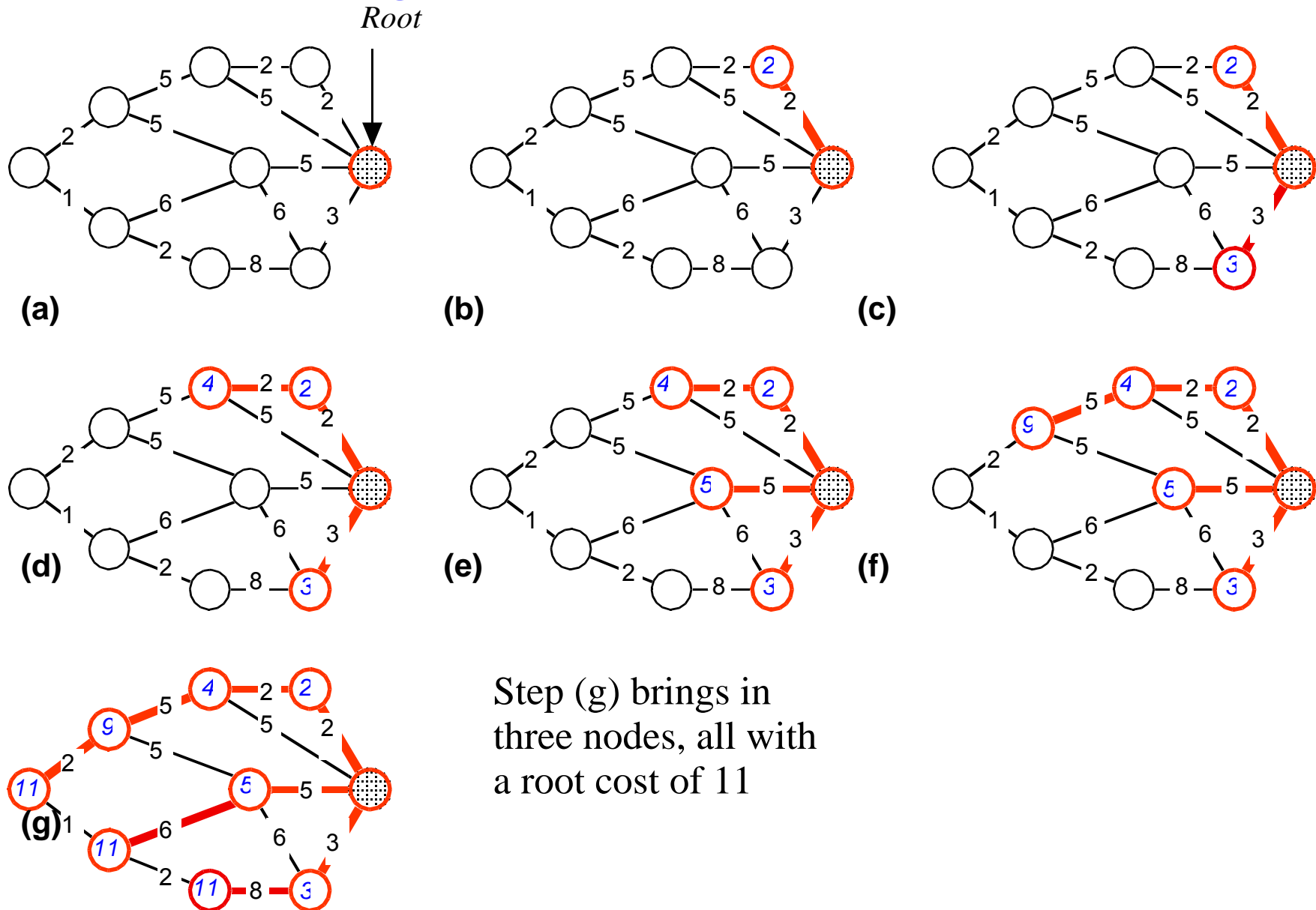
Dijkstra: the computation

- Initialize each distance $D(R1, Rn) = \infty$, except $D(R1, R1) = 0$
- Mark all D's as *tentative*, except $D(R1, R1)$ as *final*
- For each arc $Rx \rightarrow Ry$ in the graph (starting with $R1$)
 if $D(R1, Ry) > D(R1, Rx) + W(x, y)$
 then $D(R1, Ry) := D(R1, Rx) + W(x, y)$
- When all paths to Ry have been checked, Ry is removed from scanning of graph and its current $D(R1, Ry)$ is marked as *final*
 - i.e. we only minimise the path to Ry *once*; in Bellman-Ford we did it for every path through Ry
- $W(x, y) < 0$ is not allowed

Dijkstra paths (snapshot)



CS314 traditional Dijkstra example: progressive evaluation



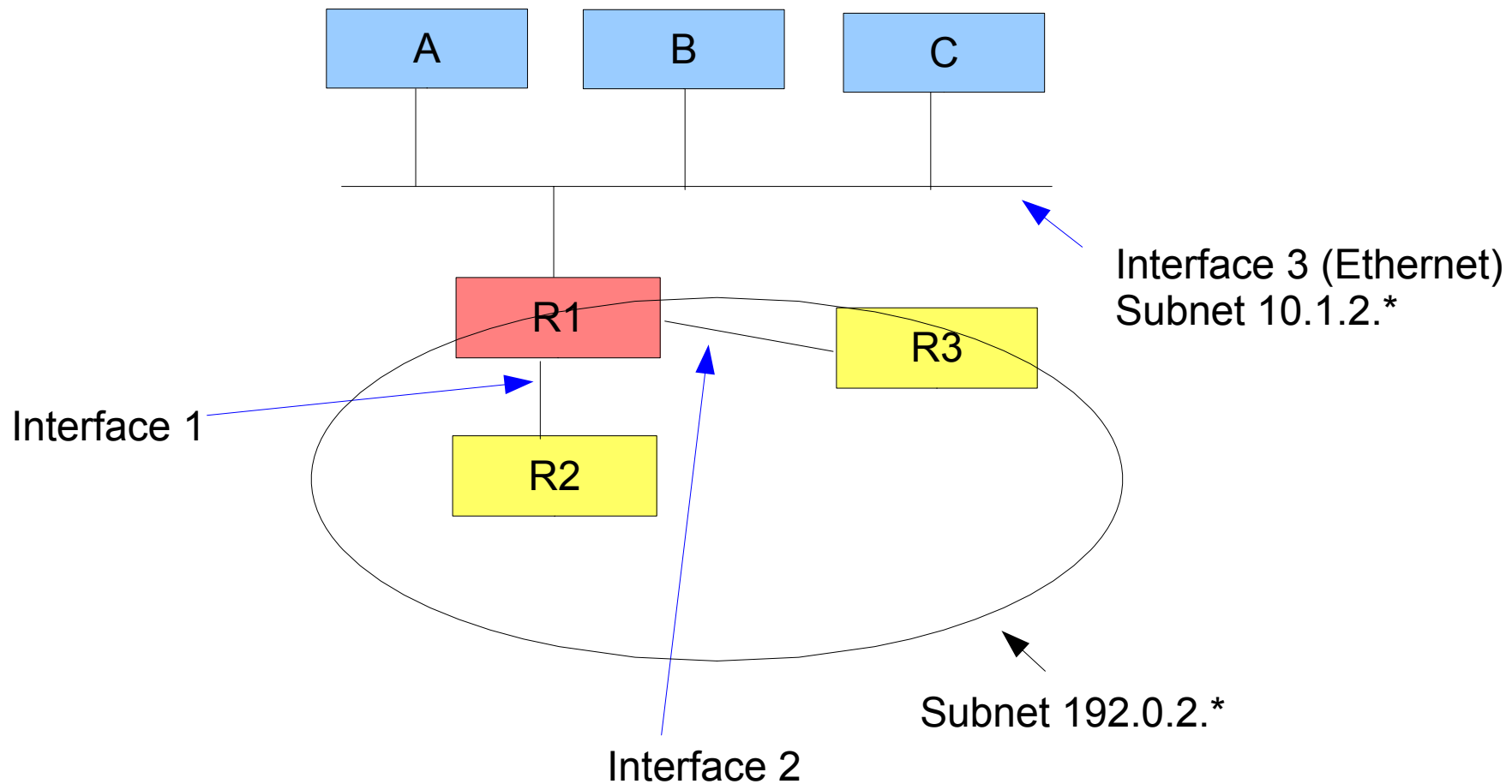
Link State: Comments

- Algorithm processes many fewer arcs than Bellman-Ford – once a distance is *final* that node is no longer considered
 - Whether this saving matters depends on size of graph and frequency of changes
- Also, each node receives neighbours' link state tables so doesn't have to start from $D(R1, Rn) = \infty$
 - Trust neighbours' path lengths – OK if you can trust your neighbours
- Link State versus Distance Vector was quite a battle in the technical community some years ago, so we have a variety of routing protocols in service today. We'll talk about three: RIP (DV), OSPF (LS) and BGP (modified DV).

RIP: Routing Information Protocol

- The original Internet DV protocol for on-site use
 - Remember that in real life, Bellman-Ford is executed as a distributed algorithm – each router computes the paths from itself to others.
 - Thus, each router keeps a table of destinations:
 - IP address of destination
 - IP address of first hop towards destination
 - Interface number for first hop
 - Distance metric for destination
 - Timestamp when entry was last updated
 - The router's complete distance vector can be read out of this table

R1 and its interfaces



Notional Initial DV table for R1

Name	Dest IP	1 st hop IP	Interface	Distance	Timestamp
R1	192.0.2.1	0.0.0.0	0	0	0
R2	192.0.2.2	0.0.0.0	1	∞	0
R3	192.0.2.3	0.0.0.0	2	∞	0
A	10.1.2.3	0.0.0.0	3	∞	0
B	10.1.2.5	0.0.0.0	3	∞	0
C	10.1.2.17	0.0.0.0	3	∞	0

Notes:

- R1 points to itself using the loopback address and distance zero.
- The other distances are set to infinity.
- R1 also shows itself as the 1st hop to its neighbours.
- All timestamps are initialized before we start.
- In practice, end systems like A, B, C don't run RIP; the router can discover them on the LAN and is pre-configured with the summarised destination....

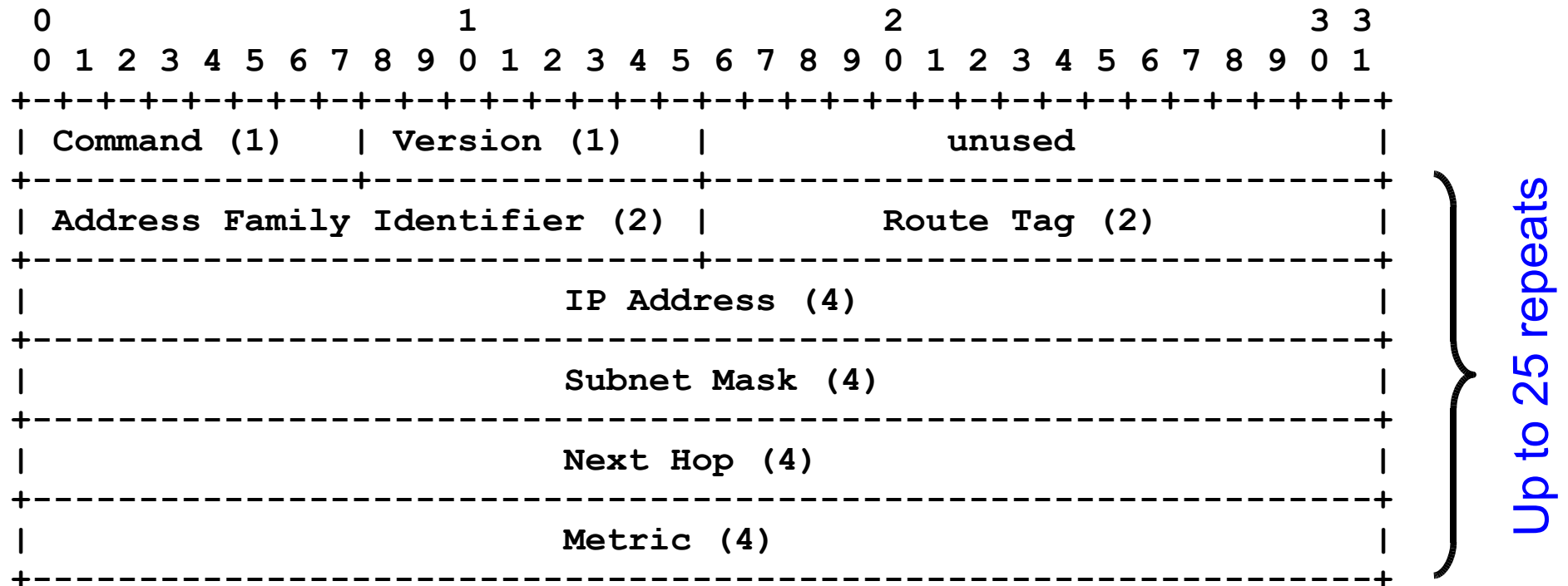
Actual Initial DV table for R1

Name	Dest IP	1 st hop IP	Interface	Distance	Timestamp
R1	192.0.2.1	0.0.0.0	0	0	0
R2	192.0.2.2	0.0.0.0	1	16	0
R3	192.0.2.3	0.0.0.0	2	16	0
None	10.1.2.*	0.0.0.0	3	1	0

Notes:

- There is no way out of configuring this externally; the router cannot guess any of it.
- In reality, destinations are address + bit mask; A, B and C can be summarised as address 10.1.2.0, mask 255.255.255.0 . This means all addresses from 10.1.2.0 to 10.1.2.255 .
- 10.1.2.* stands for 10.1.2.0, mask 255.255.255.0 .
- RIP uses 16 as infinity
- Now we have a simple distance vector that can be passed to R2 and R3.

RIP messages from router to router



Get used to this form of protocol data unit drawing – it's how all basic Internet standards show bits and bytes. This shows 24 bytes in 6 rows of 4 bytes.

RIP protocol fields

- Command
 - 1: request sending of neighbour's DV
 - 2: response including sender's DV
- Version
 - Version of RIP in use – should be 2
- AFI
 - 2 for IPv4
- Route tag
 - A value that distinguishes routes within the RIP domain from routes that go outside the RIP domain (e.g. wide area routes)
- Address and mask – as above
- Next hop address – if not the current router
 - Zero means the default route via current router
- Metric
 - A value up to 15 (number of hops)

Why is infinity equal to 16?

- By convention, RIP considers a metric of 16 to be infinity.
- Real networks should never be designed so that any normal path needs 16 hops
 - Even my messy example has at most 8 hops
- By having a low value for “infinity”, RIP will not waste too many cycles when re-computing routes after a link breakage
 - Even if a destination has become unreachable, we will stop looking for a path when the length exceeds 15

Constructing a RIP packet for R1 (1)

Command: (2)	00000010			
Version: (2)	00000010			
Unused:	00000000	00000000		
AFI,tag: (2+2)	00000000	00000010	00000000	00000000
R1's address:	11000000	00000000	00000010	00000001
Mask: (none)	00000000	00000000	00000000	00000000
Next hop:(self)	00000000	00000000	00000000	00000000
Metric: (0)	00000000	00000000	00000000	00000000
AFI,tag: (2+2)	00000000	00000010	00000000	00000000
R2's address:	11000000	00000000	00000010	00000010
Mask: (none)	00000000	00000000	00000000	00000000
Next hop:(self)	00000000	00000000	00000000	00000000
Metric: (16)	00000000	00000000	00000000	00010000

Constructing a RIP packet for R1 (2)

AFI,tag: (2+2) 00000000 00000010 00000000 00000000

R3's address: 11000000 00000000 00000010 00000011

Mask: (none) 00000000 00000000 00000000 00000000

Next hop:(self) 00000000 00000000 00000000 00000000

Metric: (16) 00000000 00000000 00000000 00010000

AFI,tag: (2+2) 00000000 00000010 00000000 00000000

LAN address: 00001010 00000001 00000010 00000000

Mask: 11111111 11111111 11111111 00000000

Next hop:(self) 00000000 00000000 00000000 00000000

Metric: (1) 00000000 00000000 00000000 00000001

How RIP operates

- Each router sends its full DV in one or more Response messages
 - On startup
 - Every 30 seconds
 - Whenever something changes (e.g. a link goes down)
 - When a Request message is received
- Recalculates its DV whenever it receives a Response
 - Which includes running Bellman-Ford algorithm
- That is enough for all routers to have up-to-date routing tables
 - Unless rate of changes is too great for the messages and calculations to keep up

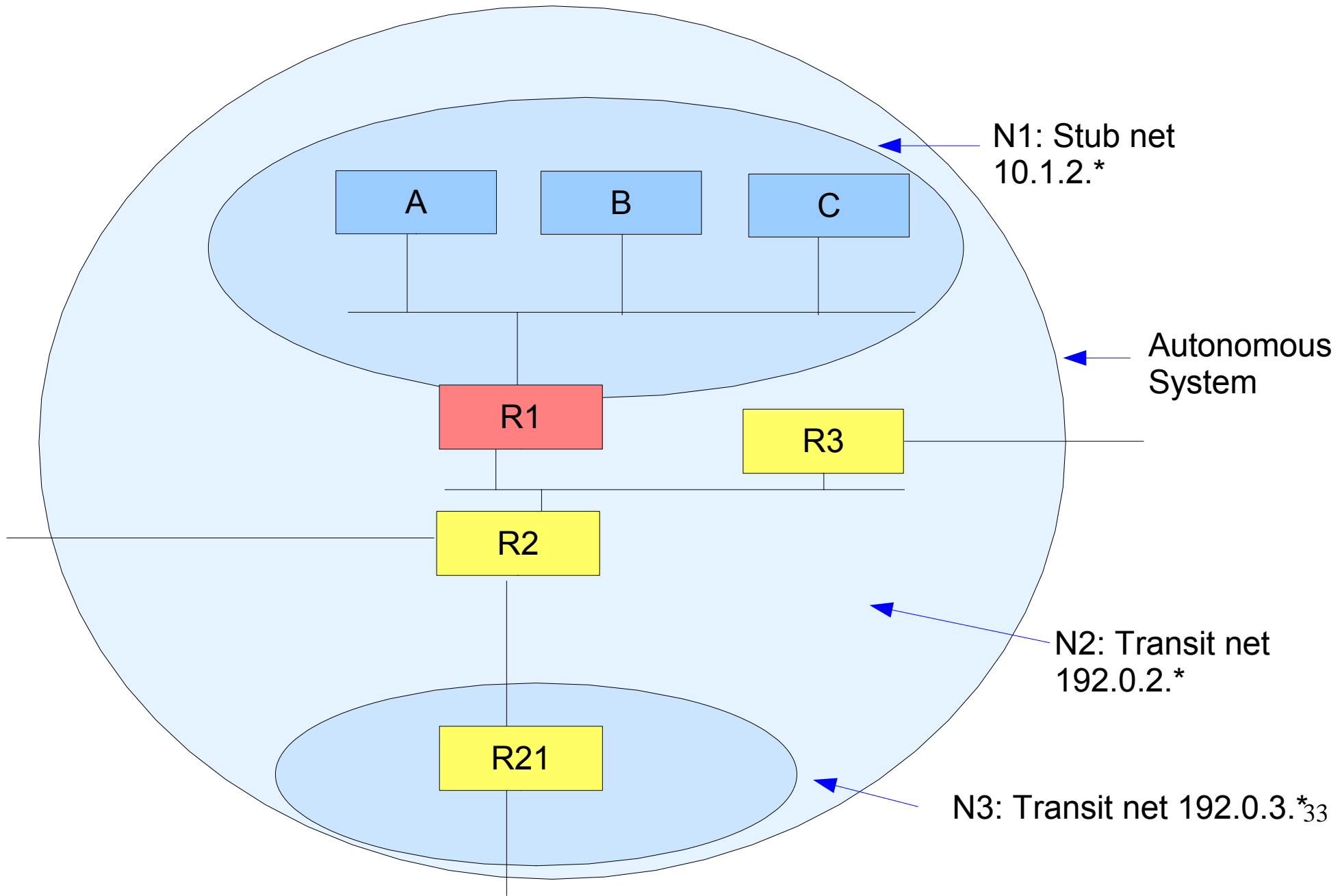
Comments on RIP

- We've skipped some details.
- Obviously, RIP messages are quite noisy with many unused bits.
- Convergence time can be large when conditions are bad (e.g. recovering from power failure in a building)
- RIP therefore doesn't work well in large, complex networks.
- Generally, OSPF is considered better these days.

OSPF: Open Shortest Path First

- A common Internet LS protocol for on-site use
 - Dijkstra is executed as a distributed algorithm – each router computes its own shortest paths.
 - Conceptually, OSPF runs within a finite domain known as an Autonomous System (AS).
 - It distinguishes routers from networks within the AS.
 - Networks can be stub or transit networks.
 - Each router keeps a Link State database,
 - The router's Link State Advertisements (LSAs) to its neighbours can be read out of this database.

OSPF conceptual model



Contents of Link State database

- The vertices in the network graph will be:
N1, N2, N3, R1, R2, R3, R21
- Arcs will be **$R_x \rightarrow R_y$, $N_x \rightarrow R_y$ or $R_x \rightarrow N_y$**
(never **$N_x \rightarrow N_y$**)
- For each destination, LS database lists
 - IP address (and mask) of destination
 - IP address of first hop towards destination
 - Interface number for first hop
 - Distance metric for destination
- External networks (outside the AS) will typically have large metrics

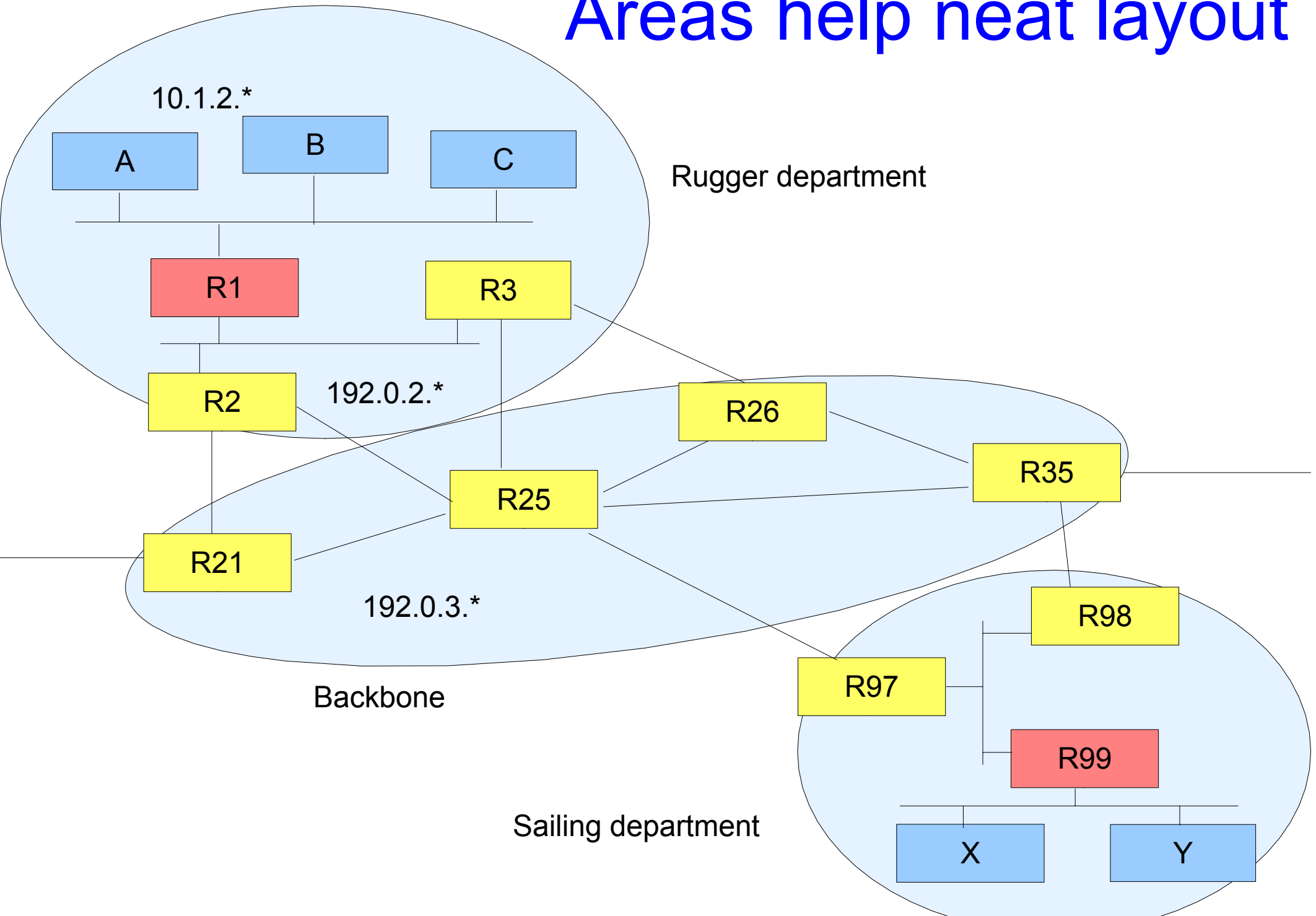
Splitting very large AS's

- OSPF already allows an AS to contain logically separate networks.
- For a large campus, it is possible to split the AS into logically separate “areas” which each run their own SPF computations.
 - Reduces overhead within each area
 - Area 0 then acts as a campus backbone
 - Inter-area routes cross the backbone
- More efficient operation in a large network
 - Glitches in one department don't affect another

Flavours of OSPF router

- Internal
 - Connected to LANs and links within an Area
- Designated
 - When >1 routers on a LAN, one is designated to send LSAs
- Area Border
 - Connected to the backbone Area and at least one normal Area. Computes routes for each.
- Backbone
 - Connected to the backbone Area
- AS Boundary
 - Connected to the outside world
 - Advertises routes to external networks

Areas help neat layout



OSPF message types

- | | | |
|---|-----------------------------|---|
| 1 | <u>Hello</u> | Discover/maintain neighbours |
| 2 | <u>Database Description</u> | Summarise database contents |
| 3 | <u>Link State Request</u> | Request database download |
| 4 | <u>Link State Update</u> | Send database update |
| 5 | <u>Link State Ack</u> | Acknowledge an update
(<u>not</u> abbreviated as LSA) |

- Link State Updates carry the LSAs (Link State Advertisements) of the sending router and LSAs from its upstream routers.
- Within an Area, LSAs describe link state in detail, similar to RIP
- Beyond an Area, LSAs summarise routes per network ³⁸

What causes an LSA to be updated?

- Its lifetime expires
- Interface up/down
- Designated router for a LAN changes
- Change of state of neighbour router
- Intra-area route changes in border router
- Inter-area route changes
- Border router attaches to new area
- Backbone topology change (virtual route change)
- External route change
- AS Boundary router down

What's in an OSPF LSA

- Five types of LSA:
 - 1 Router-LSAs – link state for each interface
 - 2 Network-LSAs – routers for a transit subnet
 - 3,4 Summary-LSAs – inter-Area link state
 - 5 AS-external-LSAs – external link state

Main items in a Router-LSA

LS type = 1 ;indicates router-LSA Link State

Link State ID = 192.0.2.1 ;R1's Router ID

Advertising Router = 192.0.2.1 ;R1's Router ID

#links = 2

Link ID = 192.0.2.3 ;IP address of Desig. Rtr.

Link Data = 192.0.2.1 ;R1's IP interface to net

Type = 2 ;connects to transit net

metric = 1

Link ID = 10.1.2.0 ;IP Network number

Link Data = 255.255.255.0 ;Network mask

Type = 3 ;connects to stub net

metric = 2

Main items in a Network-LSA

LS type = 2 ;indicates network-LSA Link State

Link State ID = 192.0.2.3 ;IP address of Desig. Rtr.

Advertising Router = 192.0.2.3 ;R3's Router ID

Network Mask = 255.255.255.0

Attached Router = 192.0.2.3 ;Router ID

Attached Router = 192.0.2.1 ;Router ID

Attached Router = 192.0.2.2 ;Router ID

How OSPF operates

- Each router sends its LSAs in one or more Link State Update messages
 - When timer expires
 - Whenever an LSA changes
 - When a LS Request message is received
- Recalculates its LSAs whenever it receives an LSU Response
 - Which includes running Dijkstra algorithm
- That is enough for all routers to have up-to-date routing tables
- We've skipped many details.
 - Convergence times and scalability significantly better than RIP

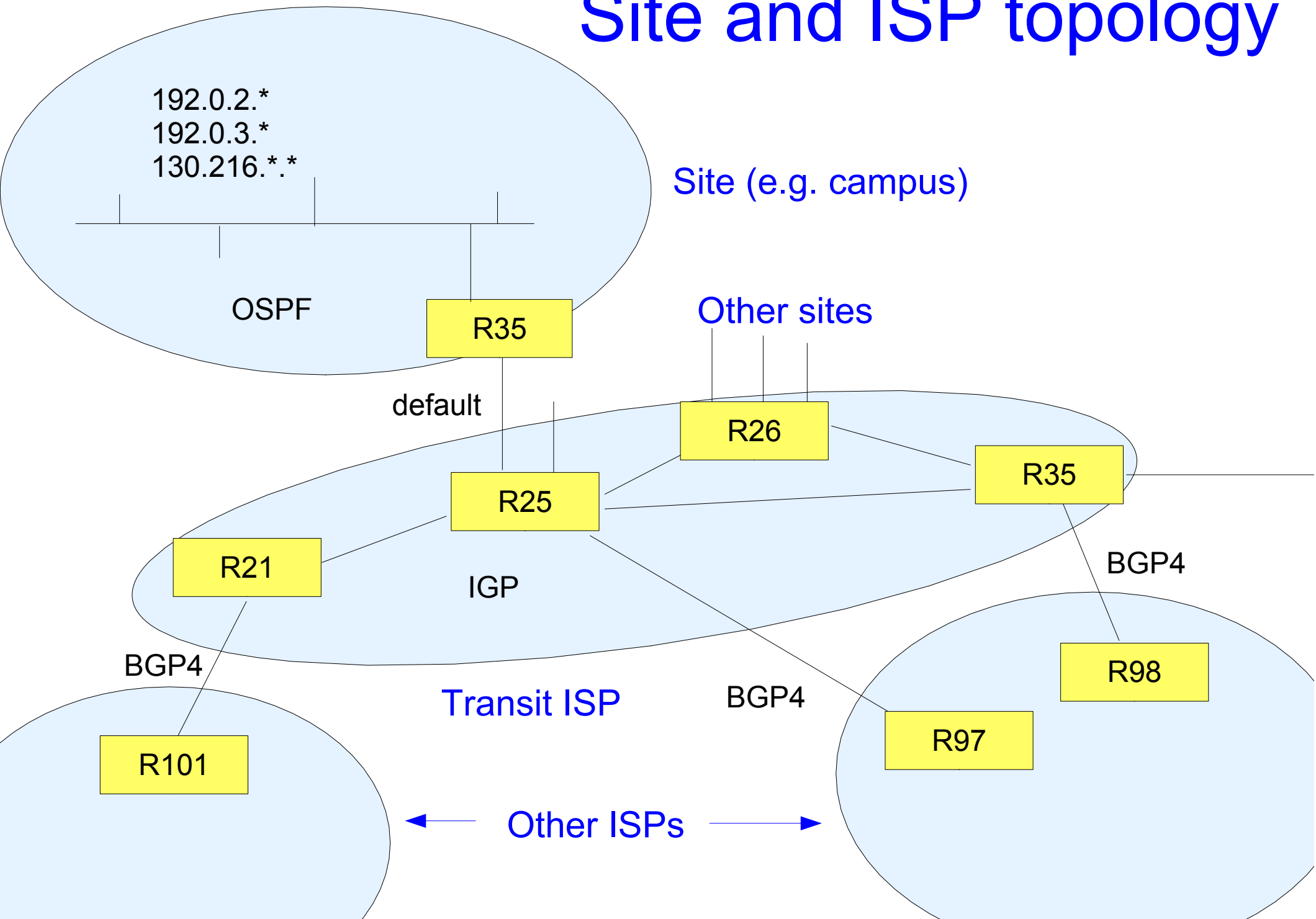
BGP4: Border Gateway Protocol v4

- The Internet distinguishes two main classes of routing protocols
 - Interior Gateway Protocols (IGPs) used within an organisation or campus, e.g., RIP or OSPF
 - Exterior Gateway Protocols (EGPs) used between organisations or campuses, e.g. BGP4
 - BGP4 has no competition today.
 - EGPs are also sometimes called Inter-Domain Routing (IDR) protocols
- BGP4 is a modified Distance Vector protocol, sometimes called a Path Vector protocol

The problem BGP solves

- Assume every site is running an IGP (OSPF or an alternative). The world now needs to find routes to the site, for a set of networks, e.g. 192.0.2.*, 192.0.3.*, 130.216.*.*
 - A site may have several unrelated address blocks.
 - A consumer ISP looks like a site
 - In the world, there are hundreds of 1000s of sites
- How can any site find routes to any other site?
 - Starting point: the *default route* 0.0.0.0 in the IGP points to a site exit router
 - The ISP accepts packets via an ingress router
 - It needs IGP routes to its customer sites and BGP routes₄₅ to everything else

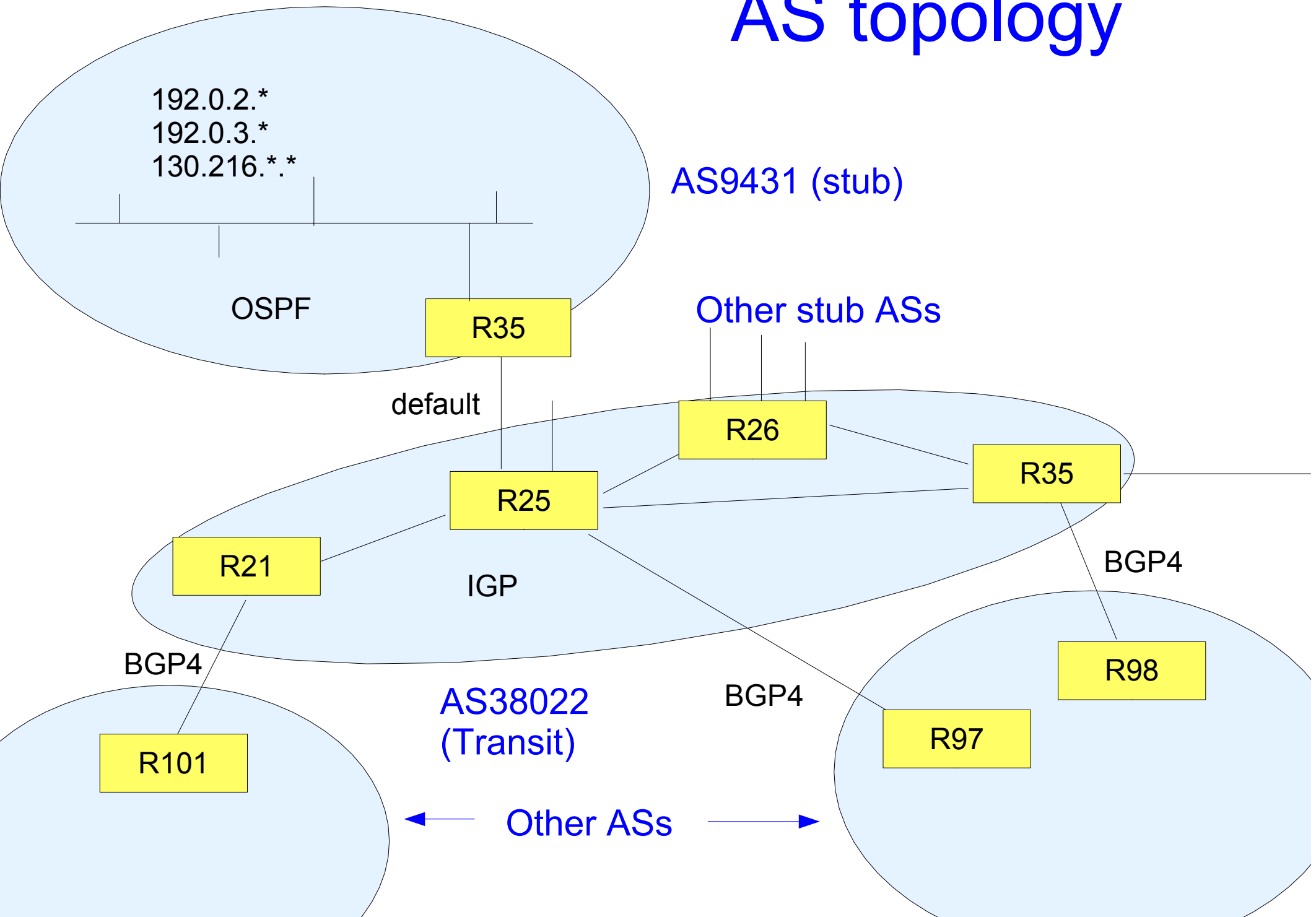
Site and ISP topology



BGP4 AS Paths

- BGP formalises the concept of Autonomous System
 - Over-simplifying, each ISP and each major customer has an AS number.
 - Several ASs can in fact be federated so they appear externally as a single AS.
- An AS path is basically a list of AS numbers that leads to a given destination
 - A destination is a network address/mask, usually called a *prefix* in BGP
 - 130.216/16 means a 16 bit prefix: 130.216.0.0 with mask 255.255.0.0

AS topology



About AS numbers

- An AS path to 130.216/16 will end with ..., AS38022, AS9431
- AS numbers are currently 16 bits but will be expanded to 32 bits, since >43,000 have been handed out already.

BGP messages

- Open
 - Sent when a link to a neighbour BGP speaker comes up
- Update
 - Routing information (see below)
- Notification
 - Sent when a link to a neighbour BGP speaker is intentionally closed
- Keepalive

Note – no Request message in BGP

BGP update messages

- BGP speakers send updates to their BGP neighbours
- They contain various *attributes* per destination
 - Origin of information (IGP or EGP)
 - AS Path
 - IP address of next hop (if not the sender of the update)
 - Unreachable (if announcing a path failure)
 - Metric (only within an AS; BGP does not weight links)
 - Community (see below)

BGP communities

- BGP supports routing *policies* instead of weights.
- The *community* attribute, if used, acts as a label for a group of destinations, so that policy rules can be applied to the group.
 - Artificial example rule: do not use paths to this community that include AS5168.

BGP operation

- When a BGP router receives an update
 - It filters the update according to local policy rules, discarding AS paths that its policy rejects.
 - Then runs Bellman-Ford over the resulting DV table (but uses AS path length as the metric, hence the phrase “path vector”).
 - Then generates update messages for its other neighbours.
- When a BGP router observes a link failure
 - It updates its DV table, runs Bellman-Ford, and generates update messages for its neighbours.

Aggregation

- Today's global BGP4 table has about 230,000 entries (>5 per AS number).
 - See history at <http://bgp.potaroo.net/>
- It's only that small because BGP *aggregates* adjacent prefixes in a binary tree.
 - 192.0.2.0/24 and 192.0.3.0/24 can be aggregated as 192.0.2.0/23, etc.
 - Two prefixes that share an AS path are aggregated if possible, before advertising the path.

BGP in practice

- BGP4 is the routing protocol that makes the Internet work.
 - Without BGP aggregation, the Internet would have melted down many years ago, with several times more entries in the global table.
 - Today, hardware improvements are roughly keeping up with growth.
 - But the number of dynamic updates is some cause for concern. What if updates start to arrive faster than they can be processed and passed on to the neighbours?

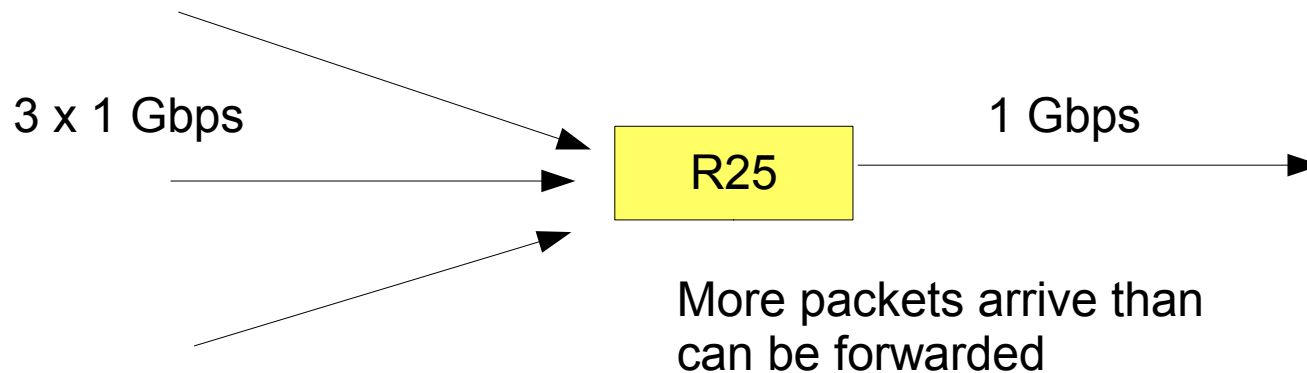
Missing topics: MPLS and IS-IS

- No time in this course to discuss MPLS (MultiProtocol Label Switching).
 - This is a “layer 2.5” technique for building semi-permanent switched paths across a backbone network.
 - Used today by many ISPs underneath the traditional IGP/EGP router layer.
- We also haven't discussed IS-IS (Intermediate System - Intermediate System).
 - An alternative to OSPF used by many ISPs.

RIBs and FIBs – building a router

- **RIB: Routing Information dataBase**
 - The routing table resulting from RIP, OSPF, BGP4, etc.
 - Generally, it's computed and accessed by software.
- **FIB: Forwarding Information dataBase**
 - A simple table that tells the hardware which interface to send a packet on:
 - IP Address – Mask – Interface number
 - Generally written by software, and read by specialist hardware operating at line speed (100 Mbps to several Gbps)
 - At 1 Gbps you have ~4 microseconds to process a packet, so the FIB lookup must take a fraction of a microsecond.

Handling congestion



- Normally, networks are designed to avoid such bottlenecks, but sometimes bursts of traffic will cause congestion and buffer overflow.
 - Buffer overflow at both ends of a link could cause deadlock, with both ends unable to send to the other.
- IP is an “unreliable” protocol - congested routers discard excess traffic unpredictably when buffers fill up.
 - We'll discuss later how this affects upper layer protocols

Routing – related to distance

- How do we connect devices

- In the same room?
- In the next room? (tens)
- In the next building? (100s)
- In the next department?
- In the next campus? (1000s)
- In the next country? (millions)
- In the next continent? (billions)

Hubs, switches

Cable/radio boundary

OSPF

+ Admin boundary

+ Funding boundary

BGP

+ Political boundary

MPLS

Distance, logical
separation, cost and
numbers increase; must
share costs and cables

References

- Shay 10.3 to 10.8
- Advanced text book (*completely optional*):

Interconnections (2nd edition)
by Radia Perlman (Addison-Wesley)
Published 1999 but still unbeaten

- The technical specs for all Internet protocols are published free as numbered “RFCs” (originally “Request for Comment”). They make tough reading and are *completely optional* for this course.

<http://www.rfc-editor.org/rfcsearch.html>

- RIP: RFC 1723
- OSPF: RFC 2328
- BGP4: RFC 4271