

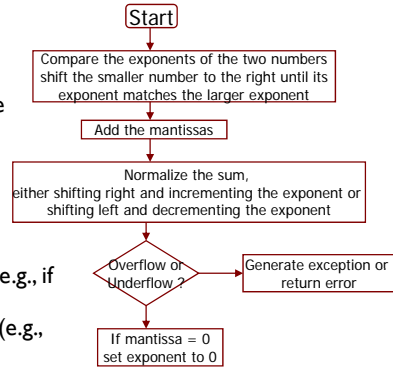
IEEE 754 Floating Point Operations

Basic Operations

- Addition: $X + Y = (M_x * 2^{Ex-E_y} + M_y) * 2^{E_y}$, $Ex \leq E_y$
- Subtraction: $X - Y = (M_x * 2^{Ex-E_y} - M_y) * 2^{E_y}$, $Ex \leq E_y$

Procedures for addition/subtraction:

- Adjust exponents and align mantissa
 - The exponent of the operands must be made equal for addition and subtraction.
 - If $E_y > E_x$ Right shift M_x to form $M_x * 2^{Ex-Ey}$
 - If $Ex > E_y$ Right shift M_y to form $M_y * 2^{E_y-Ex}$
- Add or subtract mantissa
- Normalize the result
 - Left shift result, decrement result exponent (e.g., if result is 0.001xx...) or
 - Right shift result, increment result exponent (e.g., if result is 10.1xx...)
- Check result
 - Overflow/underflow
 - If result mantissa is 0, may need to set the exponent to zero to return a zero.



Flow chart for add operation

Floating Point Addition

• $1.25 + 0.25$

◦ 0 01111111 010...000 + 0 01111101 000...000

Sign bit

Exponent bits

Mantissa bits

Steps:

- Adjust exponents and align mantissa
 - Start by adjusting the smaller exponent to be equal to the larger exponent
 - Take 0.25 (0 01111101 000...000) (with smaller exponent)
 - Original Value: E:01111101 M:000...000
 - Shifted 1 place: E:01111110 M:100...000 (Note: "1" is the hidden bit)
 - Shifted 2 places: E:01111111 M:010...000
- Add mantissa bits
 - 0 01111111 1.010...000
 - +0 01111111 0.010...000
 - 0 01111111 1.100...000
- Normalize result:
 - No need to change. It is normalized.
- Check result
 - OK. Answer = 0 01111111 100..000

The hidden bit

Floating Point Subtraction

- **1.25 - 0.25**
 - $0\ 01111111\ 010\dots000 - 0\ 01111101\ 000\dots000$
 - Steps:
 - Adjust exponents and align mantissa
 - Start by adjusting the smaller exponent to be equal to the larger exponent
 - Take 0.25 ($0\ 01111101\ 000\dots000$) (with smaller exponent)
 - Shifted 1 place: E:01111110 M:100...000 (Note: "1" is the hidden bit)
 - Shifted 2 places: E:01111111 M:010...000
 - Subtract mantissa bits
 - $0\ 01111111\ 1.010\dots000$
 - $\underline{-0\ 01111111\ 0.010\dots000}$
 - $0\ 01111111\ 1.000\dots000$
 - Normalize result:
 - No need to change. It is normalized.
 - Check result
 - OK. Answer = $0\ 01111111\ 000\dots000$

Multiplication

- **Floating Point Operations:**
 - $X = (-1)^{S_x} M_x * 2^{E_x}$, $Y = (-1)^{S_y} M_y * 2^{E_y}$
 - Multiplicatoin: $X * Y = (M_x * M_y) * 2^{E_x+E_y}$
- **Procedures for Multiplicatoin:**
 - Check Zeros
 - If one or both operands is equal to zero, return the result as zero.
 - Compute the sign of the result $S_x \text{ XOR } S_y$
 - Multiply mantissa
 - $M_x * M_y$
 - Round the result to the allowed number of mantissa bits
 - Add exponents
 - biased exponent (E_x) + biased exponent (E_y) - bias
 - Normalize the result
 - Left shift result, decrement result exponent (e.g., 0.001xx...)
 - Right shift result, increment result exponent (e.g., 10.1xx...)
 - Check result
 - If larger/smaller than maximum exponent allowed return exponent overflow/underflow

Floating Point Multiplication

- **-18 * 9.5**
 - | 1000011 0010...000 * 0 1000010 0011...000
 - Steps:
 - Sign = 0 XOR 1 = 1
 - Multiply mantissa (don't forget the hidden bit)
 - $$\begin{array}{r} 1.0010 \\ * 1.0011 \\ \hline 10010 \\ 10010 \\ 00000 \\ 00000 \\ \hline 10010 \\ 101010110 = 1.01010110 \end{array}$$
 - Add exponents
 - $1000\ 0011 + 1000\ 0010 - 01111111 = 1000\ 0110$
 - Normalize the result
 - It is normalized. No change
 - Check result
 - OK. Answer = 1 10000110 01010110...0

Division

- **Floating Point Operations:**
 - $X = (-1)^{S_x} M_x * 2^{E_x}$, $Y = (-1)^{S_y} M_y * 2^{E_y}$
 - Division: $X / Y = (M_x / M_y) * 2^{E_x - E_y}$
- **Procedures for Division:**
 - Check Zeros
 - If both operands is equal to zero, return the result as NaN
 - If Y is equal to zero, return the result as infinity.
 - Compute the sign of the result S_x XOR S_y
 - Divide mantissa
 - M_x / M_y (Round the result to the allowed nbr of mantissa bits)
 - Subtract exponents
 - Division: biased exponent (E_x) - biased exponent (E_y) + bias
 - Normalize the result
 - Left shift result, decrement result exponent (e.g., 0.001xx...)
 - Right shift result, increment result exponent (e.g., 10.1xx...)
 - Check result
 - If larger/smaller than maximum exponent allowed return exponent overflow/underflow

Floating Point Division

- 3.75 / 1.5
 - 0 10000000 1110...000 / 0 01111111 100...000
 - Steps:
 - Sign = 0 XOR 0 = 0
 - Divide mantissa (don't forget the hidden bit)
 - $$\begin{array}{r} \underline{1\ 01} \\ 11 \overline{) 11.11} \\ \underline{11} \\ 0 = \end{array}$$
 - = 1.01
 - Subtract exponents
 - $10000000 - 01111111 + 01111111 = 1000\ 0000$
 - Normalize the result
 - It is normalized. No change
 - Check result
 - OK. Answer = 0 10000000 0100..000

Conversion Examples:

- Example 1
 - What is the IEEE floating point representation of 6.5_{10} ? Hence, what is the representation for 52?
 - Step 1: 6.5_{10} :
 - Sign = 0
 - $6.5_{10} \Rightarrow$ Binary number = 110.1
 - Normalization (shift radix point to left by 2 places) $\Rightarrow 1.101 \times 2^2$
 - Exponent = $127+2=129=10000001$
 - Mantissa = 1010...0
 - Answer = 0 10000001 1010...0 = 40D00000
 - Step 2: $52/6.5 = 8 = 2^3$ i.e. $6.5 \times 2^3 = 52$
 - Sign bit : unchanged
 - Mantissa : unchanged
 - Exponent : exponent from step 1 + 3 = $(129 + 3) = 10000100$
 - Answer = 0 10000100 1010...0 = 42500000
- Example 2
 - What is the IEEE floating point representation of 0.625_{10} ? Hence, what is the representation for 0.875_{10} ?
 - Step 1: $0.625 = 3F200000 = 0\ 01111110\ 010...0$
 - Step 2: 0.875
 - $0.875 = 0.625 + 0.25$
 - $= 0.625 + 0.5 * 2^{-1}$
 - $= 1.25 * 2^{-1} + 0.5 * 2^{-1}$
 - Sign bit: same
 - Exponent bit: same
 - Mantissa = + 0.5
 - Answer = 0 01111110 110...0 = 3F600000