

COMPSCI 210 S1T Computer Systems

Floating Point Numbers Arithmetic

Agenda & Reading

◆ Agenda:

- Floating Point operations
 - ◆ Addition
 - ◆ Subtraction
 - ◆ Multiplication
 - ◆ Division

Floating Point Operations

- ◆ Exponents and fraction must be handled separately & differently.
- ◆ Consider the Floating point numbers:
 - $X = M_x * 2^{E_x}$, $Y = M_y * 2^{E_y}$
- ◆ Basic Operations
 - Addition: $X + Y = (M_x * 2^{E_x - E_y} + M_y) * 2^{E_y}$, $E_x \leq E_y$
 - Subtraction: $X - Y = (M_x * 2^{E_x - E_y} - M_y) * 2^{E_y}$, $E_x \leq E_y$
- ◆ Procedures for addition/subtraction:
 - Adjust exponents and align mantissa
 - ◆ The exponent of the operands must be made equal for addition and subtraction.
 - ◆ If $E_y > E_x$ Right shift M_x to form $M_x * 2^{E_y - E_x}$
 - ◆ If $E_x > E_y$ Right shift M_y to form $M_y * 2^{E_x - E_y}$
 - Add or subtract mantissa
 - Normalize the result
 - ◆ Left shift result, decrement result exponent (e.g., if result is 0.001xx...) or
 - ◆ Right shift result, increment result exponent (e.g., if result is 10.1xx...)
 - Check result
 - ◆ If larger than maximum exponent allowed return exponent overflow
 - ◆ If smaller than minimum exponent allowed return exponent underflow
 - ◆ If result mantissa is 0, may need to set the exponent to zero to return a zero.

Example on Decimal Value

- ◆ $3.25 * 10^2 + 2.63 * 10^{-1}$
- ◆ Steps:
 - 1) Align decimal points:
 - ◆ Start by adjusting the smaller exponent to be equal to the larger exponent, and modify the fraction accordingly.
 - ◆ Take $2.63 * 10^{-1}$
 - Original value = $2.63 * 10^{-1}$ (right-shift mantissa-> increase exponent)
 - Shift 1 place = $0.263 * 10^0$
 - Shift 2 places = $0.0263 * 10^1$
 - Shift 3 places = $0.00263 * 10^2$
 - 2) Add the numbers
 - ◆ $3.25 * 10^2$
 - ◆ $+ 0.00263 * 10^2$
 - ◆ $3.25263 * 10^2$
 - 3) Normalize the result
 - ◆ No need to change. It is already normalized.
 - 4) Check result
 - ◆ OK. Answer = $3.25263 * 10^2$

Example on Binary value

◆ $0.101 \times 2^2 + 0.111 \times 2^4$

◆ Steps:

- 1) Align decimal points:
 - ◆ Start by adjusting the smaller exponent to be equal to the larger exponent, and modify the fraction accordingly.
 - Take 0.101×2^2
 - Original value = 0.101×2^2
 - Shifted 1 place = $0.101 \times 2^2 \Rightarrow 0.0101 \times 2^3$
 - Shifted 2 places = $0.0101 \times 2^3 \Rightarrow 0.00101 \times 2^4$
- 2) Add the numbers
 - ◆ 0.111×2^4
 - ◆ $+ 0.00101 \times 2^4$
 - ◆ 1.00001×2^4
- 3) Normalize the result
 - ◆ No need to change. It is normalized.
- 4) Check result
 - ◆ OK. Answer = 1.0011×2^4

Floating Point Addition

◆ $1.25 + 0.25$

- $0.25 = 0\ 01111101\ 000...000$
- $1.25 = 0\ 01111111\ 010...000$

◆ Steps:

- Adjust exponents and align mantissa
 - ◆ Start by adjusting the smaller exponent to be equal to the larger exponent
 - ◆ Take 0.25 (0 01111101 000...000) (with smaller exponent)
 - Original Value: E:01111101 M:000...000
 - Shifted 1 place: E:01111110 M:100...000 (Note: "1" is the hidden bit)
 - Shifted 2 places: E:01111111 M:010...000
- Add mantissa bits
 - ◆ $0\ 01111111\ 1.010...000$
 - ◆ $+0\ 01111111\ 0.010...000$
 - ◆ $0\ 01111111\ 1.100...000$
- Normalize result: The hidden bit
 - ◆ No need to change. It is normalized.
- Check result
 - ◆ OK. Answer = $0\ 01111111\ 100...000$

Floating Point Addition

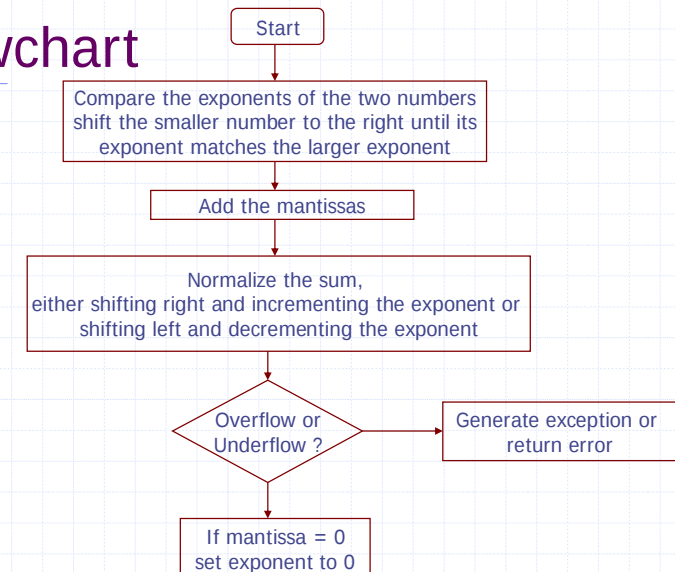
◆ $6.5 + 1.75$

- $6.5 = 0\ 10000001\ 1010...000$
- $1.75 = 0\ 01111111\ 110...000$

◆ Steps:

- Adjust exponents and align mantissa
 - ◆ Start by adjusting the smaller exponent to be equal to the larger exponent
 - ◆ Take 1.75 (0 01111111 110...000) (with smaller exponent)
 - Original Value: E:01111111 M:110...000
 - Shifted 1 place: E:10000000 M:1100...000 (Note: "1" is the hidden bit)
 - Shifted 2 places: E:10000001 M:01110...000
- Add mantissa bits
 - ◆ $0\ 10000001\ 1.1010...000$
 - ◆ $+0\ 10000001\ 0.01110...00$
 - ◆ $0\ 10000001\ 10.0001000...0$
- Normalize result:
 - ◆ Shift the mantissa right by 1 bit increase the exponent by 1
 - ◆ E: 10000001 + 1 = 10000010, mantissa = 1.0000100..0
- Check result:
 - ◆ OK. Answer = $0\ 10000010\ 0000100...000$

Flowchart



Floating Point Subtraction

◆ 1.25 - 0.25

- 0.25 = 0 01111101 000...000
- 1.25 = 0 01111111 010...000

◆ Steps:

- Adjust exponents and align mantissa
 - ◆ Start by adjusting the smaller exponent to be equal to the larger exponent
 - ◆ Take 0.25 (0 01111101 000...000) (with smaller exponent)
 - Shifted 1 place: E:01111110 M:100...000 (Note: "1" is the hidden bit)
 - Shifted 2 places: E:01111111 M:010...000
- Subtract mantissa bits
 - ◆ 0 01111111 1.010...000
 - ◆ -0 01111111 0.010...000
 - ◆ 0 01111111 1.000...000
- Normalize result:
 - ◆ No need to change. It is normalized.
- Check result
 - ◆ OK. Answer = 0 01111111 000...000

Floating Point Subtraction

◆ 6.5 - 1.75

- 6.5 = 0 10000001 1010...000
- 1.75 = 0 01111111 110...000

◆ Steps:

- Adjust exponents and align mantissa
 - ◆ Start by adjusting the smaller exponent to be equal to the larger exponent
 - ◆ Take 1.75 (0 01111111 110...000) (with smaller exponent)
 - Shifted 1 place: E:10000000 M:1110...000 (Note: "1" is the hidden bit)
 - Shifted 2 places: E:10000001 M:01110...000
- Subtract mantissa bits
 - ◆ 0 10000001 1.10100...000
 - ◆ -0 10000001 0.01110...000
 - ◆ 0 10000001 1.001100...000
- Normalize result:
 - ◆ No need to change.
- Check result:
 - ◆ OK. Answer = 0 10000001 00110...00

Addition Notes

◆ Truncation Errors

- If the exponents differ by more than 24, the smaller number will be shifted right entirely out of the mantissa field, producing a zero mantissa.
- The sum will then equal the larger number.
- Such truncation errors occur when the numbers differ by a factor of more than 2^{24} , which is approximately 1.6×10^7 .
- Thus, the precision of IEEE single precision floating point arithmetic is approximately 7 decimal digits.

◆ Alternatively, floating point subtraction is achieved simply by inverting the sign bit and performing addition of signed mantissas.

- Negative mantissas are handled by first converting to 2's complement and then performing the addition.
- After the addition is performed, the result is converted back to sign-magnitude form.

◆ When adding numbers of opposite sign, cancellation may occur, resulting in a sum which is arbitrarily small, or even zero if the numbers are equal in magnitude.

Multiplication

◆ Floating Point Operations:

- $X = (-1)^{S_x} M_x * 2^{E_x}$, $Y = (-1)^{S_y} M_y * 2^{E_y}$
- Multiplicatoin: $X * Y = (M_x * M_y) * 2^{E_x + E_y}$

◆ Procedures for Multiplicatoin:

- Check Zeros
 - ◆ If one or both operands is equal to zero, return the result as zero.
- Compute the sign of the result $S_x \text{ XOR } S_y$
- Multiply mantissa
 - ◆ $M_x * M_y$
 - ◆ Round the result to the allowed number of mantissa bits
- Add exponents
 - ◆ biased exponent (E_x) + biased exponent (E_y) - bias
- Normalize the result
 - ◆ Left shift result, decrement result exponent (e.g., if result is 0.001xx...) or
 - ◆ Right shift result, increment result exponent (e.g., if result is 10.1xx...)
- Check result
 - ◆ If larger than maximum exponent allowed return exponent overflow
 - ◆ If smaller than minimum exponent allowed return exponent underflow

Decimal Multiplication/Division

◆ Example 1: $3.0 * 10^1 * 4.5 * 10^2$

◆ Steps:

- Sign = $0 \text{ XOR } 0 = 0$
- Multiply mantissa (don't forget the hidden bit)
 - ◆ $3.0 * 4.5 = 13.5$
- Add exponents
 - ◆ $1 + 2 = 3$
- Normalize the result
 - ◆ $13.5 * 10^3$, Shift 1 place $\rightarrow$ increase exponent = $1.35 * 10^4$
- Check result
 - ◆ OK. Answer = $1.35 * 10^4$

◆ Example 2: $4.0 * 10^3 * 0.5 * 10^2$

- Sign = $0 \text{ XOR } 0 = 0$
- Divide mantissa (don't forget the hidden bit)
 - ◆ $4.0 / 0.5 = 8.0$
- Subtract exponents
 - ◆ $3 - 2 = 1$
- Normalize the result
 - ◆ $8.0 * 10^1$ No change.
- Check result
 - ◆ OK. Answer = $8.0 * 10^1$

Floating Point Multiplication

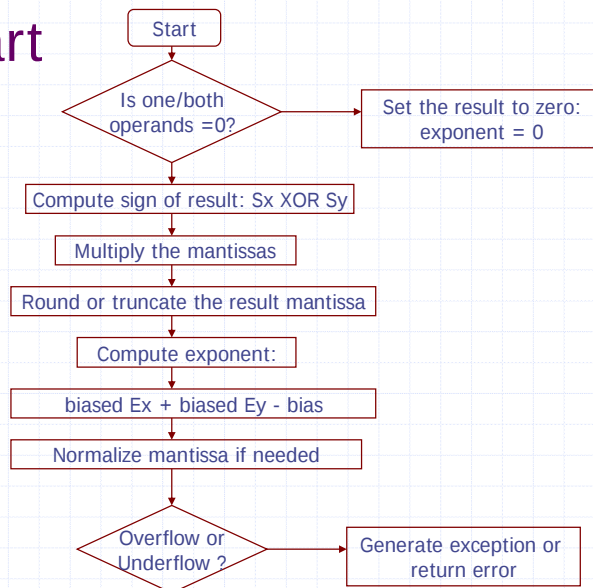
◆ -18 * 9.5

- -18 = 1 10000011 0010...000
- 9.5 = 0 10000010 0011...000

◆ Steps:

- Sign = $0 \text{ XOR } 1 = 1$
- Multiply mantissa (don't forget the hidden bit)
 - ◆ 1.0010
 - ◆ $* 1.0011$
 - ◆ 10010
 - ◆ 10010
 - ◆ 00000
 - ◆ 00000
 - ◆ 10010
 - ◆ $101010110 = 1.01010110$
- Add exponents
 - ◆ $1000\ 0011 + 1000\ 0010 - 01111111 = 1000\ 0110$
- Normalize the result
 - ◆ It is normalized. No change
- Check result
 - ◆ OK. Answer = 1 10000110 01010110...0

Flowchart



Division

◆ Floating Point Operations:

- $X = (-1)^{S_x} M_x * 2^{E_x}$, $Y = (-1)^{S_y} M_y * 2^{E_y}$
- Division: $X / Y = (M_x / M_y) * 2^{E_x - E_y}$

◆ Procedures for Division:

- Check Zeros
 - ◆ If both operands is equal to zero, return the result as NaN
 - ◆ If Y is equal to zero, return the result as infinity.
- Compute the sign of the result $S_x \text{ XOR } S_y$
- Divide mantissa
 - ◆ M_x / M_y
 - ◆ Round the result to the allowed number of mantissa bits
- Subtract exponents
 - ◆ Division: biased exponent (E_x) - biased exponent (E_y) + bias
- Normalize the result
 - ◆ Left shift result, decrement result exponent (e.g., if result is 0.001xx...) or
 - ◆ Right shift result, increment result exponent (e.g., if result is 10.1xx...)
- Check result
 - ◆ If larger than maximum exponent allowed return exponent overflow
 - ◆ If smaller than minimum exponent allowed return exponent underflow

Floating Point Division

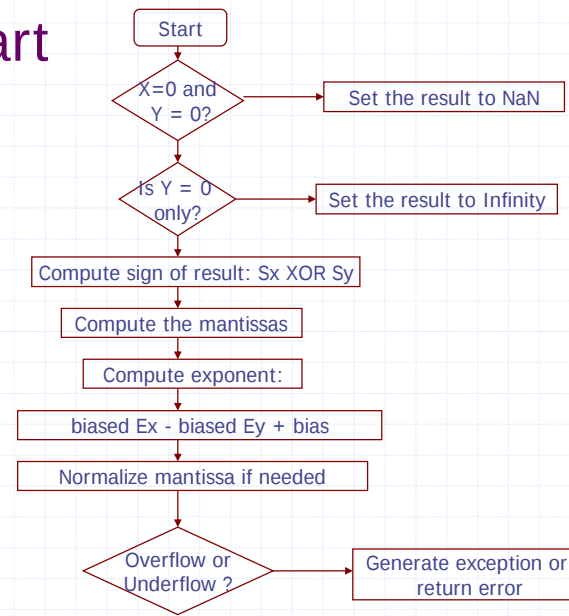
◆ 3.75 / 1.5

- 3.75 = 0 10000000 1110...000
- 1.5 = 0 01111111 100...000

◆ Steps:

- Sign = 0 XOR 0 = 0
- Divide mantissa (don't forget the hidden bit)
 - ◆
$$\begin{array}{r} 101 \\ 11 \overline{)1111} \\ \underline{11} \\ 0 \end{array}$$
 - ◆ = 101 => 1.01
- Subtract exponents
 - ◆ $10000000 - 01111111 + 01111111 = 1000\ 0000$
- Normalize the result
 - ◆ It is normalized. No change
- Check result
 - ◆ OK. Answer = 0 10000000 0100...000

Flowchart



Conversion Example:

◆ What is the IEEE floating point representation of 6.5_{10} ? Hence, what is the representation for 52?

- Step 1: 6.5_{10} :
 - ◆ Sign = 0
 - ◆ $6.5_{10} \Rightarrow$ Binary number = 110.1
 - ◆ Normalization (shift radix point to left by 2 places) $\Rightarrow 1.101 \times 2^2$
 - ◆ Exponent = $127+2=129=10000001$
 - ◆ Mantissa = 1010...0
 - ◆ Answer = 0 10000001 1010...0 = 40D00000
- Step 2: $52/6.5 = 8 = 2^3$ i.e. $6.5 \times 2^3 = 52$
 - ◆ Sign bit : unchanged
 - ◆ Mantissa : unchanged
 - ◆ Exponent : exponent from step 1 + 3 = $(129 + 3) = 10000100$
 - ◆ Answer = 0 10000100 1010...0 = 42500000

Conversion Example (con't)

◆ What is the IEEE floating point representation of 1.25_{10} ? Hence, what is the representation for 1.75 and 1.875?

- Step 1: $1.25 = 3FA00000 = 0\ 01111111\ 010...0$
- Step 2: 1.75
 - ◆ $1.75 = 1.25 \times 2^0 + 0.5 \times 2^0$
 - ◆ Sign bit: unchanged
 - ◆ Exponent bit: unchanged
 - ◆ Mantissa = + 0.5
 - ◆ Answer = 0 01111111 110...0 = 3FE00000
- Step 3:
 - ◆ $1.875 = 1.25 \times 2^0 + 0.5 \times 2^0 + 0.125 \times 2^0$
 - ◆ Sign bit: unchanged
 - ◆ Exponent bit: unchanged
 - ◆ Mantissa = + 0.5 + 0.125
 - ◆ Answer = 0 01111111 1110...0 = 3FF00000

◆ Note: exponent = $2^0 = 1$ only

- If exponent is not equal to 1, you need to take this into account

Conversion Example (con't)

- ◆ What is the IEEE floating point representation of 0.625_{10} ?
Hence, what is the representation for 0.875_{10} ?

- Step 1: $0.625 = 3F200000 = 0\ 01111110\ 010...0$

- Step 2: 0.875

- ◆ $0.875 = 0.625 + 0.25$
- ◆ $= 0.625 + 0.5 * 2^{-1}$
- ◆ $= 1.25 * 2^{-1} + 0.5 * 2^{-1}$
- ◆ Sign bit: same
- ◆ Exponent bit: same
- ◆ Mantissa = $+ 0.5$
- ◆ Answer = $0\ 01111110\ 110...0 = 3F600000$