Scheduling

- To maximize the use of the computing machinery the OS needed to know:
- what and how many devices a process would use
- what and how many files a process would use
- how much output it was likely to produce
- how long it was expected to take
- Some of this information was submitted on the control cards, e.g. the files to be used. Some of this information was inferred by the job queue the job was submitted on. Different queues meant different expected resource requirements. Jobs on different queues were scheduled differently.

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Power to the people

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Circa 1970s.

Hardware was cheaper. This meant two things:

- · Computer systems could afford to be nicer to users.
- People were now the expensive bit, we wanted to make them more productive.

Give them all a console

- CRT VDUs, early versions used teletypewriters.
- The early ones still looked like printers.
- Text was displayed on the screen top to bottom / left to right.
- You couldn't easily edit in place.
- Flushing the whole screen at 2400 baud was a bit obvious.
- Side effect people do different things with computers
- Shock, horror! Word processing, email



Scheduling (cont.)

Queue	1	2	3	4
Processor time (in seconds) default	120	20	20	120
maximum	1200	120	20	2000
I/O time (in seconds) default	120	20	20	120
maximum	1200	120	20	2000
Printer (in lines) default	500	500	300	500
maximum	2400	1200	500	3600

This system was used at the University of Auckland.

- The user could specify limits up to the maximum allowed for a queue. If no limit was specified the default for that queue was applied.
- Queue 3 was the default. It was designed for jobs which were quite small and which would move through the system smoothly. They were not allowed to use magnetic or paper tape (hence they did not require operator intervention).
- Queue 2 jobs were allowed to use two tapes (input or output).
- Most large jobs went through Queue 1.
- Queue 4 was for very large jobs and these jobs were scheduled manually.

To stop users abusing the system, they were limited to two jobs in any one queue at a time.

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Time-sharing systems

People want good response

- waiting over 1/5 of a second is noticeable
- waiting over 5 seconds is unacceptable

People working at the machine are unpredictable

- they want to use files and devices at any time
- it is very hard for the OS to work out what processes will be compute intensive or IO intensive
- more processes are active in the system

there is usually at least one process to deal with commands for each user, plus the normal work they requested

- the terminal is used for the entering of data and programs as well as JCL type commands
- because people are slow peripherals many more had to be on-line than jobs needed to be running on a batch system to get through the same amount of work
- users would also communicate with each other early email systems and talk programs

Time-sharing systems (cont.)

Time-sharing changes

Devices and the CPU can't be used to capacity otherwise the response time is too long.

So there has to be slack in the system.

·look at the processor usage on your Windows PC using System monitor or Task Manager •or run top on UNIX

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Batch system remnants

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- Ways of running processes at specific times without user intervention e.g. crontab, at and Task Scheduler.
- Until the arrival of cheap graphics terminals the terminal still looked very like a card reader. e.g the vt52.
- Command files or shell scripts very like the control cards of a job control language.

- Over-demand for resources can happen easily (but we hope occasionally).
- High-level scheduling decisions are made by the users.

Security is more of a problem

Hacking on a batch system was difficult. The system must have some authentication process.

More user administration.

Profiles, defaults, resource limits

Different types of users - System administrators with more rights than "ordinary" users.

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Desktop computers

Circa 1980s

Started with resident monitor systems.

TRS-DOS Apple DOS MS-DOS MacOS

File systems



There was no need for security. Only one user. No hard disks.

Only one program ran at a time. People usually work with one program at a time.

Gradual need for multiprogramming Print spooling TSRs Mac Switcher

Personal computers (cont.)

- Xerox (late 70s early 80s) had done some pioneering work on making computers easier to use. This was originally on time-sharing and networked workstations.
- needed high definition graphics screens
- desktop metaphor make the computer look like something else (even though *files* were called *files* well before this)
- Macintosh (now officially called "Mac") took this and popularized a cut-down version. MS-DOS users laughed scornfully until 1995.



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Networks

1980s on

New problems

security providing transparent access to resources developing protocols

Network Operating System

provides file sharing provides communication scheme runs independently from other computers on the network

Distributed Operating System

less autonomy between computers gives the impression there is a single operating system controlling the network.

Both are known as

loosely coupled as opposed to tightly coupled systems

Personal computers (cont.)

Eventually personal computers became as powerful as "real computers". Time-sharing OSs could fit on the desktop. UNIX was the main example.

Fully featured PC operating systems were developed with

virtual memory multiprogramming sophisticated file systems networking multi-user support e.g. MacOS X, Windows 8, Linux, FreeBSD

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Multiprocessor systems

Tightly coupled system – processors share memory and a clock; communication usually takes place through the shared memory.

Almost all computers now count as tightly coupled systems.

Advantages of parallel systems

Increased *throughput* Economical for the increase in performance Increased reliability graceful degradation fail-soft systems (shut down non-essentials)

Symmetric multiprocessing (SMP)

Each processor runs an identical copy of the operating system.Many processes can run at once without performance deterioration.Most modern operating systems support SMPessential now that we have multiple core CPUs

Asymmetric multiprocessing

Each processor is assigned a specific task; master processor schedules and allocates work to slave processors.

More common in extremely large systems

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Real-time systems

- A completely different thread of OS development.
- Real-time systems have varying levels of timing constraints.
- Hard real-time must satisfy requests within definite time periods or the system fails.
- e.g. robotics, air traffic control, nuclear power plants
- Soft real-time it doesn't matter too much if a time constraint is not met exactly.
- e.g. Multi-media software, telephone system
- Most modern OSs can handle some sort of soft real-time control.
- Hard real-time control systems have to be specially designed.

Pocket computers and smartphones

- PalmOS (before WebOS) small memory and slow processors. The OS had to be very efficient in order to get reasonable performance. No memory protection, no virtual memory.
- Android Based on Linux, programming applications is done in Java to Android specific APIs. A cut down version of Linux for embedded devices.
- iOS Based on MacOS X. Virtual memory, includes paging for code but not for data (why?)

SymbianOS – also true memory protection.

Other problems of size and battery consumption. GUI, amount of RAM that can be kept valid.

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Before the next lecture

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Read textbook sections

Chapter 16 Virtual Machines