## COMPSCI 107 Computer Science Fundamentals

Lecture 12 – ADTs and Stacks

# Software Engineering Design Principle

#### Modularity

Divide the program into smaller parts

## Advantages

- Keeps the complexity managable
- Isolates errors (parts can be tested independently)
- Can replace parts easily
- Eliminates redundancies (each part takes care of own data)

Easier to write, Easier to read, Easier to test, easier to modify

# Hide implementation within the module



# Abstraction

## Separates the purpose from the implementation

#### Procedural abstraction

- Specification (function header and return type) separate from implementation
- Can replace implementation if required

## Data Abstraction

Think about what can be done with the data, separate from how it is done

## Example:

- Mapping of keys to values
- Use two parallel lists
- Use a dictionary

## A collection of data and a set of operations on that data

- Specifications of an ADT focus on what the operations are
- Implementation is not specified

## ADT is not a data structure

- Data structure is a construct within a programming language
- List, tuple, dictionary

# ADT

## ADT

 An interface which is visible to the user of the ADT (the client )



## Data structure

A data structure used to implement the ADT

A wall of ADT operations isolates the data structure and the implementation from the program that uses it.

## **ADT - Integers**

#### Data

[..., -3, -2, -1, 0, 1, 2, 3, ...]

## Operations

- Addition
- Subtraction
- ...
- Equality
- Ordering
- Representation for printing

## ADT - Set

#### Data

An unordered collection of unique elements

## Operations

- Add
- Remove
- Union
- Intersection
- Complement

#### Data collection

- Position of the elements matters and is stable
- Two ends to the structure (front and back, first and last)

Different structures has different ways to add and remove elements

# ADT - Stack

#### Ordered collection of data

- Addition of items and removal of items happens at the same end
- Top of the stack

## Remove data in reverse order of data added

Last in first out (LIFO)

## Operations

- Push
- Pop
- Peek
- Is\_empty
- Size

- Implementation using Python list
- What is the big-O of push()?
- What is the big-O of pop()?

```
class StackV1:
```

```
def __init__(self):
    self.items = []
```

```
def is_empty(self):
    return self.items == []
```

def push(self, item):
 self.items.insert(0,item)

def pop(self): return self.items.pop(0)

def size(self): return len(self.items)

- Implementation using Python list
- What is the big-O of push()?
- What is the big-O of pop()?

```
class StackV2:
def __init__(self):
```

```
self.items = []
```

```
def is_empty(self):
    return self.items == []
```

```
def push(self, item):
    self.items.append(item)
```

```
def pop(self):
    return self.items.pop()
```

def size(self): return len(self.items)

# **Checking braces**

- Many computer languages use braces to signify start and end
  - They need to be matched to be correct
  - They need to be nested correctly
- Stacks can be used to determine correct use of braces
- Algorithm:
  - Add each open brace to the stack
  - When a closing brace is encountered, check to see if a matching brace is on the top of the stack
  - When the last token is checked, the stack should be empty

## Exercise

Use a stack to check if the braces are used correctly in the following strings. Show the state of the stack after each stack operation.

# • ( ( ) ) • ( ( ) ( • [ ( ) { } ] • [ [ ( ) ] [ { } ] ) • [ [ ( ) ] [ { } ] )

- To be done after class: Write a program that uses a Stack to check if an input string has correctly matching braces
- check\_braces('(this) [is] {best} ([done] {with} (stacks))')

- Standard mathematics represents expressions using infix notation
  - Operators appear between the operands



- Postfix notation puts the operator after the operands
  - No brackets are needed to specify order of precedence

## 4 5 +

- Convert the following infix expressions into postfix notation
- 4 \* 5 2 \* 8 1
- 1 4 + 2 \* 3 \* 5
- 9 6 / 4 + 2 \* (2 + 3)
- 2 \* ((4 + 2) \* 3 + 2) 1

# Converting from infix to postfix

- A stack can be used in the algorithm to convert infix to postfix
  - Divide expression into tokens
  - Operators: +. -, \*, /
  - Operands: single digits
  - Other tokens: brackets

# Algorithm for converting infix to postfix

- Create a stack to store operators and a list for the output tokens
- Scan the tokens from left to right
- If the token is an operand, add it to the output list
- If the token is a left parenthesis, push it to the operator stack
- If the token is a right parenthesis, pop the operator stack until the left parenthesis is removed. Append each operator to the output list
- If the token is an operator, push it onto the operator stack. But first, remove any operators that have higher or equal precedence and append them to the output list
- When there are no more tokens, remove operators on the stack and append to the output list

Show the operator stack and the output list at every step as the following infix expression is converted to postfix

- Create an empty stack
- Scan the list of tokens from left to right
- If the token is an operand, push it to the operand stack
- If the token is an operator, pop the stack twice
  - The first element popped is the right operand
  - The second element popped is the left operand
- Apply the operator to the operands and push the result onto the stack
- When there are no more tokens, the stack should contain the result.

Following the algorithm to evaluate postfix expressions, show the operand stack, and the token being processed (at each step) as the following postfix expression is evaluated:

## 7 12 8 9 - \* 3 / +